

NT-Cryptography
MAT4930 0329

Home-U

Prof. JLF King
Wedn., 15Feb2023

Due: ~~Monday, 20Feb.~~ no later than 11:30AM, on Wedn., 22Feb., slid completely under my office door, LIT402.

Fill-in every blank on this sheet. This sheet is the first-page of your write-up.

U1: Alice publishes her ElGamal modUlus $U := 4094957$, gen. $G := 399510$, and her public key $A := \langle G^\alpha \rangle = 859311$, where α is Alice's private key, and $\langle \cdot \rangle$ means $\langle \cdot \rangle_U$. Bob transmits his public key $B := \langle G^\beta \rangle = 856746$. Each computes $\sigma = \langle G^{\alpha\beta} \rangle$, the secret key. Bob skipped class on known plaintext day, and erroneously ElGamal's messages m_0, \dots, m_9 to Alice, but reusing β . He transmits

$C_0 := 2501615$ $C_1 := 1685151$ $C_2 := 20561$ $C_3 := 2079233$
 $C_4 := 2287623$ $C_5 := 2428749$ $C_6 := 990351$ $C_7 := 3630623$
 $C_8 := 39151$ $C_9 := 1225900$; ten ciphertexts $C_j := \langle \sigma \cdot m_j \rangle$.

Eve knows Bob sent his [crummy] password, $M_K := 11111$, and she tricked him into sending $M_C := 4930$, their Crypto course number. Bob's error, together with the Known and Chosen plaintexts, allow you, Eve, to compute $\sigma =$ and recover all ten plaintexts. Eve used what property of M_C that M_K might not possess?

For b -bit modulus U , with Bob sending N messages [one known, one chosen plaintext], what is the running time $R(b, N)$ of Eve's algorithm to compute σ ?

U2: RSA uses a modulus N , (en/de)cryption exponents E, d so that $E \cdot d = 1 + k\varphi(N)$, for some posint k . In class, we restricted Bob's message m to be $\perp N$, then used EFT to conclude that $m^{Ed} \equiv_N m$.

Pair (m, N) is nice if: $\forall k \in \mathbb{N}: m^{1+k\varphi(N)} \equiv_N m$. Posint N is great if (m, N) is nice for every integer m .

i Prove that each $N := pq$, with $p < q$ primes, is great.

ii Characterize, with proof, the set of great numbers.

U3: i Use Pollard- ρ to find a nt-factor of $M := 59749$, using seed $s_0 := 7$ and map $f(x) := \langle 1+x^2 \rangle_M$. Make a nice table, labeled

Time | Tortoise | Hare | $s_{2k} - s_k$ | GCD(??)

—but replace the “??” with the correct expression. You found non-trivial factor $E :=$

The hare Hits into the tortoise at time $H :=$

Repeat, showing the table for $s_0 := 24$. Experiment with different seeds; what is the typical running time? [RT means $\#(f\text{-evals})$]. How is it related to the factor you find?

ii A seed s determines a tail; the smallest natnum T for which there is a time $n > T$ with $f^n(s) = f^T(s)$. The smallest such n is $T+L$ where L is the period. Derive (picture+reasoning) a formula for the hitting time $H(T, L)$. [Hint: $H(0, L) = L$.]

iii Produce a Floyd-like algorithm that computes both T and L . The number, N , of f -evaluations is upper-bounded by some small constant times $T+L$ (=arclength of ρ). How small can you get $N(T, L)$? [Hint: When $T = 0$, Floyd's Tortoise-Hare alg. uses $3L$ evaluations.] Your Floyd-like alg. may be able to upper-bnd the $f\text{-eval}$ # in form $\alpha T + \beta L$, for specific posints α, β . [Is $T=0$ a special case?]

U4: Bob's RSA modulus is $M := p \cdot q$, where $p < q$ are b -bit primes. Doofusly, Bob wrote value $F := \varphi(M)$ on a paper napkin, which Eve found. Describe Eve's algorithm to rapidly compute p in time $O(b^n)$, where $n =$ $\in \mathbb{Z}_+$.

[Assume, for every k -bit target T , that sqroot,remainder $s, r \in \mathbb{N}$ satisfying $[s^2] + r = T < [s+1]^2$, can be found in $O(k^2)$ time.]

End of Home-U

U1: 115pts

U2: 115pts

U3: 85pts

U4: 35pts

Total: 350pts

HONOR CODE: “I have neither requested nor received help on this exam other than from my team-mates and my professor (or his colleague).” Name/Signature/Ord

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