

Fibonacci sequences

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See also *Problems/Algebra/LinearAlg/linear-recurr.latex*

GCD. We'll sometimes use a red under-bracket for greatest common divisor, e.g., $\underline{n, m} := \text{GCD}(n, m)$.

Prolegomenon. The famous *Fibonacci sequence* $\vec{f} := (\mathbf{f}_n)_{n=-\infty}^{\infty}$ is defined by $\mathbf{f}_0 := 0$, $\mathbf{f}_1 := 1$ and

$$1a: \quad \mathbf{f}_{n+1} = \mathbf{f}_n + \mathbf{f}_{n-1},$$

producing this doubly- ∞ sequence:

$$\begin{aligned} n \dots & -5 \ -4 \ -3 \ -2 \ -1 \ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \dots \\ \mathbf{f}_n \dots & 5 \ -3 \ 2 \ -1 \ 1 \ 0 \ 1 \ 1 \ 2 \ 3 \ 5 \ 8 \ 13 \ 21 \ 34 \ 55 \dots \end{aligned}$$

Let α and β be the positive and negative roots of the *characteristic polynomial* of \vec{f} , which is

$$\text{Fib}(x) := x^2 - x - 1 \stackrel{\text{note}}{=} [x - \alpha][x - \beta]. \text{ So}$$

$$1b: \quad \alpha + \beta = 1 \quad \text{and} \quad \alpha \cdot \beta = -1. \text{ Moreover} \\ \alpha^2 = \alpha + 1, \beta^2 = \beta + 1 \text{ and } \alpha, \beta = \frac{1}{2}[1 \pm \sqrt{5}].$$

For future reference,

$$1c: \quad \alpha > 1 > |\beta|.$$

Binet Formula. Let $\varphi := \sqrt{5}$. [Mnemonic “phi” to evoke “five”.] Our \vec{f} is some linear combination $A \cdot [n \mapsto \alpha^n] + B \cdot [n \mapsto \beta^n]$. Easily, $-B = A = \frac{1}{\varphi}$, so

$$2: \quad \forall n \in \mathbb{Z}: \quad \mathbf{f}_n = \frac{1}{\varphi} \cdot [\alpha^n - \beta^n],$$

since this formula gives correct values for \mathbf{f}_0 and \mathbf{f}_1 .

2a: **Lemma.** *Courtesy Binet:*

$$\dagger: \quad \forall n \in \mathbb{Z}: \quad \alpha^n = \mathbf{f}_n \alpha + \mathbf{f}_{n-1}. \quad \text{Also,}$$

$$\ddagger: \quad 5\alpha^2 = [\alpha + 2]^2.$$

Each of these holds with α replaced by β . ◊

Pf of (\dagger) . We prove (\dagger) for n positive. The base case is $\alpha^1 = 1 \cdot \alpha + 0 = \mathbf{f}_1 \cdot \alpha + \mathbf{f}_0$. Inductively

$$\begin{aligned} \alpha^{n+1} &= \alpha \cdot [\mathbf{f}_n \alpha + \mathbf{f}_{n-1}] = \mathbf{f}_n \alpha^2 + \mathbf{f}_{n-1} \alpha \\ &= \mathbf{f}_n [\alpha + 1] + \mathbf{f}_{n-1} \alpha \\ &= [\mathbf{f}_n + \mathbf{f}_{n-1}] \alpha + \mathbf{f}_n \end{aligned}$$

which equals $\mathbf{f}_{n+1} \alpha + \mathbf{f}_n$, as desired. ◆

Pf (\ddagger) . Squaring, $[\alpha + 2]^2 = \alpha^2 + 4[\alpha + 1] = \alpha^2 + 4\alpha^2$. ◆

Recurrence doubling

A fibonacci-like sequence $\vec{z} := (z_n)_{n=0}^{\infty}$ is defined via (possibly complex) numbers \mathcal{P} and \mathcal{S} , by

$$3.1: \quad z_{n+2} := \mathcal{S}z_{n+1} - \mathcal{P}z_n,$$

and some initial condition (z_1, z_0) . With

$$3.2: \quad \mathbf{G} := \begin{bmatrix} \mathcal{S} & -\mathcal{P} \\ 1 & 0 \end{bmatrix}, \quad \text{then } \begin{bmatrix} z_{n+1} \\ z_n \end{bmatrix} = \mathbf{G}^n \cdot \begin{bmatrix} z_1 \\ z_0 \end{bmatrix},$$

for each integer n .

A number $r \neq 0$ engenders seq $n \mapsto r^n$. It satisfies recurrence (3.1) IFF r is a root of polynomial

$$f(x) := x^2 - \mathcal{S}x + \mathcal{P} \xrightarrow{\text{factored}} [x - \mu][x - \nu].$$

Equating coeffs in the polynomial gives:

$$3.3: \quad \begin{aligned} \mu + \nu &= \mathcal{S}. & (\text{Sum}) \\ \mu \cdot \nu &= \mathcal{P}. & (\text{Product}) \end{aligned}$$

$$3.4: \quad \mu^2 = \mathcal{S}\mu - \mathcal{P} \quad \text{and} \quad \nu^2 = \mathcal{S}\nu - \mathcal{P}.$$

Henceforth, we require $\mu \neq \nu$, i.e. $\mathcal{S}^2 \neq 4\mathcal{P}$; this, since $\text{Discr}(f) \xrightarrow{\text{note}} \mathcal{S}^2 - 4\mathcal{P}$. We want $\mu \neq \nu$ so that every seq satisfying (3.1) has its n^{th} -term equal a linear combination of μ^n and ν^n .

Doubling. Sequences $(z_e)_{e \text{ even}}$ and $(z_d)_{d \text{ odd}}$ will satisfy *some* two-term linear recurrence. [The same recurrence.] We seek numbers $\hat{\mathcal{S}}$ and $\hat{\mathcal{P}}$ such that for r equaling either μ^2 or ν^2 ,

$$r^2 = \hat{\mathcal{S}}r^1 - \hat{\mathcal{P}}r^0.$$

I.e, that polynomial

$$x^2 - \hat{\mathcal{S}}x + \hat{\mathcal{P}} \xrightarrow{\text{factors}} [x - \mu^2] \cdot [x - \nu^2].$$

Squaring eqn $\mu^2 = \mathcal{S}\mu - \mathcal{P}$ gives

$$\mu^4 = \mathcal{S}^2\mu^2 - 2\mathcal{P}\mathcal{S}\mu + \mathcal{P}^2.$$

As $\mathcal{S}\mu = \mu^2 + \mathcal{P}$, so $2\mathcal{P}\mathcal{S}\mu = 2\mathcal{P}\mu^2 + 2\mathcal{P}^2$. Thus

$$\mu^4 = \underbrace{[\mathcal{S}^2 - 2\mathcal{P}]\mu^2}_{\hat{\mathcal{S}}} - \underbrace{\mathcal{P}^2}_{\hat{\mathcal{P}}}.$$

3.5: Doubling thm. Sequence $\vec{z} := (z_n)_{n=0}^{\infty}$ satisfies recurrence $z_{n+2} := \mathcal{S}z_{n+1} - \mathcal{P}z_n$, where $\mathcal{S}, \mathcal{P} \in \mathbb{C}$.

Then sequence $(z_{2n})_{n \in \mathbb{N}}$ and $(z_{2n+1})_{n \in \mathbb{N}}$ each satisfy recurrence

$$\square_{n+2} = \hat{\mathcal{S}} \cdot \square_{n+2} - \hat{\mathcal{P}} \cdot \square_n$$

where $\hat{\mathcal{S}} := \mathcal{S}^2 - 2\mathcal{P}$ and $\hat{\mathcal{P}} := \mathcal{P}^2$.

[If $\mathcal{P} \neq 0$, then \vec{z} can be extended backwards to $\vec{z} := (z_n)_{n=-\infty}^{\infty}$, and the doubling result holds.] \diamond

Proof. The foregoing argument used $\mathcal{S}^2 \neq 4\mathcal{P}$. Since sequence-values vary continuously with \mathcal{S} and \mathcal{P} , it suffices to obtain an $\mathcal{S}^2 = 4\mathcal{P}$ pair as a limit of pairs where the inequality holds. \spadesuit

Fib example. The Fib-seq has $\mathcal{S} = 1$ and $\mathcal{P} = -1$. Hence $\hat{\mathcal{S}} = 1^2 - 2 = 3$, and $\hat{\mathcal{P}} = [-1]^2$. So $(f_{2n})_{n \in \mathbb{Z}}$ and $(f_{2n+1})_{n \in \mathbb{Z}}$ each satisfy $z_{n+2} = 3z_{n+1} - z_n$.

Starting an any index, then taking every 4th-term, gives a seq satisfying $z_{n+2} = 7z_{n+1} - z_n$, since $3^2 - 2 \cdot 1 = 7$ and $1^2 = 1$. \square

3.6: Speed-up thm. Sequence $\vec{z} := (z_n)_{n=0}^{\infty}$ satisfies recurrence $z_{n+2} := \mathcal{S}z_{n+1} - \mathcal{P}z_n$.

For positint K and natnum b , subsequence $(z_{nK+b})_{n=0}^{\infty}$ satisfies $\square_{n+2} = \hat{\mathcal{S}} \cdot \square_{n+2} - \hat{\mathcal{P}} \cdot \square_n$, where $x^2 - \hat{\mathcal{S}}x + \hat{\mathcal{P}}$ is the characteristic polynomial of K^{th} -power $\mathbf{G}^K = \begin{bmatrix} \mathcal{S} & -\mathcal{P} \\ 1 & 0 \end{bmatrix}^K$. \diamond

Pf. Use $x^2 - \hat{\mathcal{S}}x + \hat{\mathcal{P}}$ for the charpoly of $\mathbf{M} := \mathbf{G}^K$. Cayley-Hamilton asserts $\mathbf{M}^2 = \hat{\mathcal{S}} \cdot \mathbf{M} - \hat{\mathcal{P}} \cdot \mathbf{I}$. Applying both sides to colvec $\begin{bmatrix} z_1 \\ z_0 \end{bmatrix}$, the bottom entry asserts

$$z_{2K} = \hat{\mathcal{S}} \cdot z_K - \hat{\mathcal{P}} \cdot z_0.$$

4: Theorem. For each integer N :

$$4a: \quad [f_N]^2 + [f_{N-1}]^2 = f_{2N-1}. \quad \diamond$$

Inefficient proof. Always, $\text{LhS}(4a)$ is non-negative. And $\text{RhS}(4a)$ is non-neg, even when $N \in \mathbb{Z}_-$, since f_{Odd} is always non-neg. So ISTProve that the squares of $\text{LhS}(4a)$ and $\text{RhS}(4a)$ are equal. To this end, define

$$4b: \quad \begin{aligned} \mathcal{L} &:= \varphi^2 \cdot [[f_N]^2 + [f_{N-1}]^2] \quad \text{and} \\ \mathcal{R} &:= \varphi^2 \cdot [f_{2N-1}] \stackrel{\text{note}}{=} \varphi \cdot [\alpha^{2N-1} - \beta^{2N-1}]. \end{aligned}$$

Leftside. By (2), $\varphi \cdot f_N = \alpha^N - \beta^N$. So $\varphi^2 \cdot f_N^2$ equals

$$\alpha^{2N} + \beta^{2N} - 2[\alpha\beta]^N.$$

But $\alpha\beta = -1$, and N and $N-1$ have opposite parities. Thus \mathcal{L} equals

$$\begin{aligned} \alpha^{2N} + \beta^{2N} + \alpha^{2[N-1]} + \beta^{2[N-1]} \\ = \alpha^{2[N-1]}[\alpha^2 + 1] + \dots, \end{aligned}$$

where the “ \dots ” represents a copy of all the α -terms to its left, but with “ α ” replaced by “ β ”.

By (1b), note, $\alpha^2 + 1 = \alpha + 2$. Thus

$$\mathcal{L} = \alpha^{2[N-1]}[\alpha + 2] + \beta^{2[N-1]}[\beta + 2].$$

Squaring \mathcal{L} will give twice this cross-term:

$$\begin{aligned} [\alpha\beta]^{2[N-1]} \cdot [\alpha + 2][\beta + 2] &= 1 \cdot [\alpha + 2][\beta + 2] \\ &= \alpha\beta + 2[\alpha + \beta] + 4 \\ &= -1 + 2 + 4 = 5. \end{aligned}$$

Also note $[\alpha + 2]^2 = \alpha^2 + 4[\alpha + 1] = 5\alpha^2$. Thus

$$\mathcal{L}^2 = \alpha^{4[N-1]} \cdot 5\alpha^2 + \dots + 5 \cdot 2.$$

Consequently

$$4c: \quad \frac{1}{5} \cdot \mathcal{L}^2 = \alpha^{4N-2} + \beta^{4N-2} + 2.$$

Rightside of (4a). Square \mathcal{R} and divide by 5. Since $5 = \varphi^2$,

$$\frac{1}{5} \cdot \mathcal{R}^2 = [\alpha^{2N-1} - \beta^{2N-1}]^2.$$

The cross-term is $-2[\alpha\beta]^{2N-1} = -2 \cdot [-1]^{2N-1} = 2$, since $2N-1$ is odd. We have thus shown that

$$4d: \quad \frac{1}{5} \cdot \mathcal{R}^2 = \alpha^{4N-2} + \beta^{4N-2} + 2.$$

And this equals $\text{RhS}(4c)$, as desired. \diamond

Ahem. It's a proof, but the next is a *prettier proof*. \square

Dot-product proof

Fix three integers $S = j + k$. Evidently the dot-product

$$\begin{aligned} f_{j+1}f_k + f_jf_{k-1} &= f_{j+1}[f_{k-1} + f_{k-2}] + f_jf_{k-1} \\ &= [f_{j+1} + f_j]f_{k-1} + f_{j+1}f_{k-2} \\ &= f_{j+2}f_{k-1} + f_{j+1}f_{k-2}. \end{aligned}$$

This last dot-product is the same as the first, but with “ $j + k = S$ ” replaced by “[$j+1$] + [$k-1$] = S ”. Thus, for all $j \in \mathbb{Z}$, expression $f_{j+1}f_k + f_jf_{k-1}$ depends only on S .

5: Theorem. For all triples of integers $S = j + k$:

$$5a: \quad f_{j+1} \cdot f_k + f_j \cdot f_{k-1} = f_S.$$

In particular, (4a) holds. \diamond

Proof. Setting $j := 0$ in $\text{LhS}(5a)$ results in

$$f_1f_S + f_0f_{S-1} = 1 \cdot f_S + 0 \cdot f_{S-1} = f_S.$$

Hence (5a). To obtain (4a) from (5a), set $S := 2N-1$ and $j := N-1$. \diamond

Alt Pf. Looking ahead to matrix $\mathbf{A} := \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$ from Cassini, the $(1, 2)$ -entry of matrix $\mathbf{A}^S = \mathbf{A}^j \mathbf{A}^k$ is (*). \diamond

5b: Duplication identity. For each integer n ,

$$*: \quad f_{2n} = f_n \cdot [f_{n+1} + f_{n-1}] = f_n \cdot [f_n + 2f_{n-1}]. \quad \diamond$$

Proof. Apply (5a) with $j := n-1$ and $k := n+1$. Then $f_{2n} = f_n f_{n+1} + f_{n-1} f_n = f_n \cdot [f_{n+1} + f_{n-1}]$. So f_{2n} equals

$$f_n \cdot [f_n + f_{n-1}] + f_{n-1} = \text{RhS}(*) . \quad \diamond$$

5c: Cassini's identity. For all n ,

$$\text{L: } f_{n+1} \cdot f_{n-1} = f_n^2 + [-1]^n. \quad \diamond$$

Pf. Let $v_n := \begin{bmatrix} f_n \\ f_{n-1} \end{bmatrix}$. With $A := \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$, our recurrence says $A v_n = v_{n+1}$. Id-matrix $I = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} f_1 & f_0 \\ f_0 & f_{-1} \end{bmatrix}$; i.e A^0 is 2×2 matrix $[v_1 \ v_0]$. Thus A^n equals $\begin{bmatrix} f_{n+1} & f_n \\ f_n & f_{n-1} \end{bmatrix}$. Equation $\text{Det}(A^n) = [\text{Det}(A)]^n$ yields (L). \diamond

Periodicity mod a prime. Working mod-5, note

$$(f_{10}, f_{11}) = (55, 89) \equiv_5 (0, -1) = [-1] \cdot (f_0, f_1).$$

Mod-5, then, Fib-seq \vec{f} has a nega-period of length 10, whence $\vec{f}_{\text{Mod-5}}$ has 20 as a period: $f_{20+n} \equiv_5 f_n$.

[Eventually adjoin material wrt other primes.]

5d: 1-3-1 Fib lemma. Fibonacci sequence \vec{f} satisfies

$$\text{Ll: } [f_{n+1}]^2 - 3[f_n]^2 + [f_{n-1}]^2 = 2 \cdot [-1]^n. \quad \diamond$$

Proof. Squaring $f_n = f_{n+1} - f_{n-1}$ implies

$$\begin{aligned} f_n^2 &= f_{n+1}^2 + f_{n-1}^2 - 2f_{n+1} \cdot f_{n-1} \\ \text{Cassini} \quad &f_{n+1}^2 + f_{n-1}^2 - 2[f_n^2 + [-1]^n] \\ &= f_{n+1}^2 + f_{n-1}^2 - 2f_n^2 - 2[-1]^n. \quad \diamond \end{aligned}$$

6: Lemma. $\forall n \in \mathbb{N}: [f_1]^2 + [f_2]^2 + \dots + [f_n]^2 = f_n f_{n+1}$. \diamond

Pf. Easy induction. Even nicer, $f_n f_{n+1}$ is the area of a $f_n \times f_{n+1}$ business card. Decompose it into squares. \diamond

Divisibility

Integer-sequence $\vec{\sigma} = (\sigma_1, \sigma_2, \dots)$ is a *divisibility sequence* [div-seq] if \forall indices $j, k \in \mathbb{Z}_+$: If $j \bullet k$ then $\sigma_j \bullet \sigma_k$.

Our $\vec{\sigma}$ is a *strong div-seq* if $\forall j, k$:

$$\dagger: \quad \sigma_{\text{GCD}(j,k)} = \text{GCD}(\sigma_j, \sigma_k).$$

Why does *strong div-seq* implies *div-seq*? When $j \bullet k$, then $\text{GCD}(j, k) = j$. Hence $\text{GCD}(\sigma_j, \sigma_k) = \sigma_j$. And this implies $\sigma_j \bullet \sigma_k$.

Strong div-seq implies, for each list k_1, \dots, k_N ,

$$\ddagger: \quad \sigma_{\text{GCD}(k_1, \dots, k_N)} = \text{GCD}(\sigma_{k_1}, \dots, \sigma_{k_N}).$$

Using underline for GCD , note $\underline{a, c, e} = \underline{a, c, e}$. Hence,

$$\begin{aligned} \underline{\sigma_{a,c,e}} &= \underline{\sigma_{a,c,e}} \xrightarrow{\text{by } (\dagger)} \underline{\sigma_{a,c}, \sigma_e} \\ &\xrightarrow{\text{by } (\dagger)} \underline{\sigma_a, \sigma_c, \sigma_e} \\ &= \underline{\sigma_a, \sigma_c, \sigma_e}. \end{aligned}$$

8.1: Lemma. For integers $b > a \geq 0$, values $\sigma_n := b^n - a^n$ form a divisibility sequence. \diamond

Proof. Given $j \bullet k$, write $k = \rho \cdot j$, and set $B := b^j$ and $A := a^j$. So

$$\frac{\sigma_k}{\sigma_j} = \frac{B^\rho - A^\rho}{B - A} = \sum_{\substack{v+u=\rho-1 \\ v,u \geq 0}} B^v A^u. \quad \diamond$$

8.2: Obs. Div-seq $\sigma_n = b^n - a^n$ need not be a strong div-sequence. E.g, with $b := 4$ and $a := 2$, note

$$\text{GCD}(\sigma_6, \sigma_4) = \text{GCD}(4032, 240) = 48.$$

Yet $\sigma_{\text{GCD}(6,4)} = \sigma_2 = 12$. \square

8.3: ?Fact? For integers $b > a \geq 1$, if $b \perp a$ then sequence $\sigma_n := b^n - a^n$ is a strong divisibility-sequence. \diamond

Proof. ? See proof of (11). \diamond

9: Theorem. Fibonacci seq $\vec{f} = (f_1, f_2, f_3, \dots)$ is a divisibility-sequence. \diamond

Pf. Fix posint P . For $n = 1, 2, \dots$, we prove $f_P \mid f_{nP}$. The base-case is $f_P \mid f_{1 \cdot P}$.

We apply Thm 5,

$$* : f_{j+k} = f_{j+1} \cdot f_k + f_j \cdot f_{k-1},$$

with $j := nP$ and $k := P$. Inductively, quotient $Q := f_{nP}/f_P$ is an integer. Our (*) says

$$\begin{aligned} f_{[n+1]P} &= f_{nP+1}f_P + f_{nP}f_{P-1} \\ &= f_{nP+1}f_P + Qf_P f_{P-1} \\ &= f_P \cdot [f_{nP+1} + Qf_{P-1}]. \end{aligned}$$

And $f_{nP+1} + Qf_{P-1}$ is an integer. \diamond

10.1: Prop. For integers δ, β, D : If $\delta \mid D$, then

$$\text{GCD}(\delta + \beta, D) = \text{GCD}(\beta, D). \quad [\text{Exercise}] \quad \diamond$$

10.2: Prop. For n an integer: $f_{n+1} \perp f_n$. $[\text{Exercise}] \quad \diamond$

11: Fib strong-div. Divide $u \neq 0$ into v [both integers] to get quotient and remainder, $v = qu + r$. Then

- 11a: $\text{GCD}(f_v, f_u) = \text{GCD}(f_u, f_{v-u})$;
- 11b: $\text{GCD}(f_v, f_u) = \text{GCD}(f_u, f_r)$;
- 11c: $\text{GCD}(f_v, f_u) = f_{\text{GCD}(v, u)}$.

Thus f_1, f_2, f_3, \dots is a strong divisibility-sequence. \diamond

Pf 11a. We apply Thm 5 with $j := v - u$ and $k := u$. So

$$f_v = f_{v-u+1}f_u + f_{v-u}f_{u-1}.$$

Since f_u divides $f_{v-u+1}f_u$, our Prop 10.1 says

$$\underline{f_v, f_u} = \underline{f_{v-u}f_{u-1}}, \quad f_u = \underline{f_{v-u}}, \quad f_u,$$

since $f_{u-1} \perp f_u$. [Argument works also for $u=0$.] \diamond

Pf 11b. Applying (11a) q times gives $\underline{f_v, f_u} = \underline{f_u, f_{v-qu}}$ \diamond

Pf 11c. Recall the update rule in the Euclidean algorithm (Lightning Bolt) when seeding LBolt with $r_0 := v$ and $r_1 := u$. Observe that the r of decomposition $v = qu + r$ is the r_2 of LBolt. Thus (11b) says

$$\underline{f_{r_0}, f_{r_1}} = \underline{f_{r_1}, f_{r_2}}.$$

But this is the update rule when seeded with f_{r_0} and f_{r_1} . Consequently, letting $g := \underline{v, u}$,

$$\underline{f_v, f_u} = \underline{f_g, f_0} = \underline{f_g, 0} \stackrel{\text{note}}{=} f_g. \quad \diamond$$

Can complex analysis Fib?

What is the RoC of

$$\dagger: \quad \mathcal{F}(z) := \sum_{n \geq 0} f_n z^n,$$

the *generating function* of the fibonacci sequence?

Recurrence $f_{n+2} = f_{n+1} + f_n$, and that $f_n \geq 0$ for $n \geq 0$ show that sequence \vec{f} is non-decreasing. For $n \geq 2$, the sequence is positive, so recurrence (1a) gives

$$2f_{n-2} \leq f_n \leq 2f_{n-1}. \quad \text{Hence}$$

$$\text{Const} \cdot [\sqrt{2}]^n \leq f_n \leq \text{Const} \cdot 2^n. \quad \text{Thus}$$

$$\frac{1}{\sqrt{2}} \geq \text{RoC}(\mathcal{F}) \geq \frac{1}{2}.$$

In particular, \mathcal{F} is analytic in a nbhd of zero.

Shh... In fact we know that

$$**: \quad f_n = \frac{1}{\sqrt{5}} \cdot [\alpha^n - \beta^n], \quad \text{where} \\ \alpha = \frac{1}{2}[1 + \sqrt{5}] \quad \text{and} \quad \beta = \frac{1}{2}[1 - \sqrt{5}].$$

Here, α and β are the *golden/silver ratios*, respectively. We thus have that

$$\text{Shh!}: \quad \text{RoC}(\mathcal{F}) = 1/\alpha = \frac{2}{1 + \sqrt{5}}.$$

Source. Online book **A FIRST COURSE IN COMPLEX ANALYSIS**, made freely available by Matthias Beck, Gerald Marchesi, Dennis Pixton & Lucas Sabalka, has the below problem in §10.3 of their text. \square

A formula for \mathcal{F} . [Below, I'll use \mathcal{F} to abbreviate $\mathcal{F}(z)$.] By defn,

$$z^0 \cdot \sum_{n \geq 0} f_{n+2} \cdot z^{n+2} = \mathcal{F} - f_0 \cdot z^0 - f_1 \cdot z^1 \stackrel{\text{note}}{=} \mathcal{F} - z;$$

$$z^1 \cdot \sum_{n \geq 0} f_{n+1} \cdot z^{n+1} = z \cdot [\mathcal{F} - f_0 \cdot z^0] \stackrel{\text{note}}{=} z \cdot \mathcal{F};$$

$$z^2 \cdot \sum_{n \geq 0} f_n \cdot z^n = z^2 \cdot \mathcal{F}.$$

From the top equality, we subtract the other two. The lefthand side of the result is

$$\sum_{n \geq 0} [f_{n+2} - f_{n+1} - f_n] z^{n+2} \stackrel{\text{note}}{=} 0,$$

since $f_{n+2} = f_{n+1} + f_n$. From the righthand sides, then, $0 = [1 - z - z^2] \mathcal{F} - z$. Hence

$$\ddagger: \quad \mathcal{F}(z) = \frac{z}{D(z)}, \quad \text{where} \quad D(z) := 1 - z - z^2.$$

Easily, $D(z) = z^2 \cdot \text{Fib}(\frac{1}{z}) = z^2 \cdot [\frac{1}{z} - \alpha][\frac{1}{z} - \beta]$. Rewriting, $D(z) = [\alpha z - 1][\beta z - 1]$. Thus

$$D(z) = \alpha \beta [z - \frac{1}{\alpha}][z - \frac{1}{\beta}] \stackrel{\text{note}}{=} -1 \cdot [z + \beta][z + \alpha].$$

Residues. With $H_n(z) := \frac{1}{z^n D(z)} \stackrel{\text{def}}{=} \frac{\mathcal{F}(z)}{z^{n+1}}$, note

$$\text{Res}(H_n) = \text{Res}_{z=0} \left(\frac{\mathcal{F}(z)}{z^{n+1}} \right) \stackrel{\text{PlexNotes P.30}}{=} \frac{1}{n!} \cdot \mathcal{F}^{(n)}(0) = f_n.$$

Looking ahead, $\frac{1}{\beta^n} - \frac{1}{\alpha^n} = \frac{\alpha^n - \beta^n}{[\alpha \beta]^n} = \frac{\alpha^n - \beta^n}{[-1]^n}$. So

$$\text{Y:} \quad \frac{1}{[-\beta]^n} - \frac{1}{[-\alpha]^n} = \alpha^n - \beta^n.$$

Contour integral. Let $C_r := \text{Sph}_r(0)$. Since $n \geq 0$, the degree of $z^n D(z)$ is at least 2, whence

$$\lim_{r \nearrow \infty} \frac{1}{2\pi i} \oint_{C_r} H_n = 0.$$

As the singularities of H_n are $0, -\beta, -\alpha$,

$$0 = \text{Res}(H_n, 0) + \text{Res}(H_n, -\beta) + \text{Res}(H_n, -\alpha), \text{ i.e} \\ *: \quad f_n = \text{Res}(-H_n, -\beta) + \text{Res}(-H_n, -\alpha).$$

Writing $-H_n(z) = \frac{1/[z^n(\alpha + z)]}{z - \beta}$, the CIF gives

$$\text{Res}(-H_n, -\beta) = 1/[-\beta]^n [\alpha - \beta]. \quad \text{Similarly,} \\ \text{Res}(-H_n, -\alpha) = 1/[-\alpha]^n [\beta - \alpha].$$

Their sum equals

$$\frac{1}{\alpha - \beta} \cdot \left[\frac{1}{[-\beta]^n} - \frac{1}{[-\alpha]^n} \right] \stackrel{\text{by (Y)}}{=} \frac{1}{\alpha - \beta} \cdot [\alpha^n - \beta^n].$$

Plugging this into (*), yields (**) – which we supposedly didn't know.

Tribonacci sequence

The bi-infinite **Trib sequence** is

$$* : \begin{aligned} (\mathbf{t}_0, \mathbf{t}_1, \mathbf{t}_2) &:= (0, 0, 1) \quad \text{and} \\ \mathbf{t}_{n+3} &:= \mathbf{t}_{n+2} + \mathbf{t}_{n+1} + \mathbf{t}_n, \end{aligned}$$

for $n \in \mathbb{Z}$. The resulting Trib sequence is
 $n: \dots -9 -8 -7 -6 -5 -4 -3 -2 -1 0 1 2 3 4 5 6 7 \dots$
 $\mathbf{t}_n: \dots 81 44 24 13 7 4 2 1 1 0 0 1 1 2 4 7 13 \dots$

A nz-complex ω engenders seq $n \mapsto \omega^n$, which satisfies the recurrence part of (*) IFF ω is a zero of $\text{Trb}(z) := z^3 - z^2 - z - 1$.

Zeros of $\text{Trb}()$. As $\text{Trb}(1) < 0 < \text{Trb}(2)$, value $\text{Trb}(p)$ is zero for some posreal $1 < p < 2$. The remaining Trb -zeros form a complex-conjugate pair $w \neq \bar{w}$.

To see this last, ISTShow that $\text{Trb}|_{\mathbb{R}}$ has but one zero. Its derivative equals

$$\text{Trb}'(x) = 3x^2 - 2x - 1 = 3[z-1][z + \frac{1}{3}].$$

At the $\text{Trb}()$ -critical-points, each of $\text{Trb}(1) = -2$ and $\text{Trb}(\frac{-1}{3}) = \frac{-22}{27}$ is negative. Hence $\text{Trb}(x)$ is negative for $x < p$, and is positive for $x > p$.

12: Prop'n. Sum $w + \bar{w} + p = 1$, product $w \cdot \bar{w} \cdot p = 1$.
 Also, $p > 1 > |w| = |\bar{w}| = \frac{1}{\sqrt{p}}$. \diamond

Pf. Write $z^3 - 1z^2 - z - 1 = [z - w][z - \bar{w}][z - p]$ and equate coeffs. \spadesuit

Trib GF. The generating function of Trib-seq is

$$\dagger: \quad \mathcal{T}(z) := \sum_{n \geq 0} \mathbf{t}_n z^n. \quad \square$$

12a: Lemma. The $\text{RoC}(\mathcal{T})$ is $1/|p|$. **Proof.** Exercise. \diamond

“Cheating”. Cardano's formula gives

$$\begin{aligned} p &= \frac{1}{3}[1 + S^+ + S^-], \\ w &= \frac{1}{6}[2 + \mu S^+ + \bar{\mu} S^-], \\ \bar{w} &= \frac{1}{6}[2 + \bar{\mu} S^+ + \mu S^-], \end{aligned}$$

where

$$\begin{aligned} \mu &:= -1 + i\sqrt{3}, & S^+ &:= \sqrt[3]{19 + 3\sqrt{33}}, \\ \bar{\mu} &:= -1 - i\sqrt{3}, & S^- &:= \sqrt[3]{19 - 3\sqrt{33}}. \end{aligned}$$

Values $S^+ > S^-$ are real. And $\text{Im}(w) > 0 > \text{Im}(\bar{w})$. Approximately $p \approx 1.839$ and $w \approx -0.419 + 0.606i$. *However*, we eschew this information FtTBeing. \square

A formula for \mathcal{T} . [Below, \mathcal{T} abbreviates $\mathcal{T}(z)$.] By defn,

$$\begin{aligned} z^0 \cdot \sum_{n \geq 0} \mathbf{t}_{n+3} \cdot z^{n+3} &= \mathcal{T} - \mathbf{t}_0 \cdot z^0 - \mathbf{t}_1 \cdot z^1 - \mathbf{t}_2 \cdot z^2 \stackrel{\text{note}}{=} \mathcal{T} - z^2; \\ z^1 \cdot \sum_{n \geq 0} \mathbf{t}_{n+2} \cdot z^{n+2} &= z \cdot [\mathcal{T} - \mathbf{t}_0 \cdot z^0 - \mathbf{t}_1 \cdot z^1] \stackrel{\text{note}}{=} z \cdot \mathcal{T}; \\ z^2 \cdot \sum_{n \geq 0} \mathbf{t}_{n+1} \cdot z^{n+1} &= z^2 \cdot [\mathcal{T} - \mathbf{t}_0 \cdot z^0] \stackrel{\text{note}}{=} z^2 \cdot \mathcal{T}; \\ z^3 \cdot \sum_{n \geq 0} \mathbf{t}_n \cdot z^n &= z^3 \cdot \mathcal{T}. \end{aligned}$$

From the top equality, subtracting the lower three gives this LhS,

$$\sum_{n \geq 0} [\mathbf{t}_{n+3} - \mathbf{t}_{n+2} - \mathbf{t}_{n+1} - \mathbf{t}_n] z^{n+3} \stackrel{\text{note}}{=} 0,$$

by (*). Hence $0 = [1 - z - z^2 - z^3] \mathcal{T} - z^2$, from the RhS. Consequently,

$$\ddagger: \quad \mathcal{T}(z) = \frac{z^2}{E(z)}, \text{ where } E(z) := 1 - z - z^2 - z^3.$$

Zeros of $E()$. As $E(z) = \text{Trb}(\frac{1}{z}) \cdot z^3$, the zeros of $E()$ are reciprocals $b := 1/p$, $m := 1/w$ and \bar{m} . Thus

$$E(z) = -1 \cdot [z - b][z - m][z - \bar{m}].$$

Residues. As $\text{RoC}(\mathcal{T})$ is positive, $\underset{z=0}{\text{Res}}\left(\frac{\mathcal{T}(z)}{z^{n+1}}\right) = \mathbf{t}_n$ for $n=0, 1, \dots$, courtesy (†).

Ratio $\frac{\mathcal{T}(z)}{z^{n+1}} \xrightarrow{\text{by } \ddagger} \frac{z}{z^n E(z)} =: H_n(z)$ is a rational function satisfying $\text{Deg}(\text{Denominator}) \geq \text{Deg}(\text{Numerator}) + 2$. Consequently, $\lim_{r \nearrow \infty} \int_{C_r} H_n(z) dz$ is zero.

For r large, circle C_r encloses $0, b, m, \bar{m}$, the four singularities of H_n . Hence their H_n -residues sum to zero. Recall that the $z=0$ residue is \mathbf{t}_n . So...

$$\mathbf{t}_n = \underset{z=b}{\text{Res}}(-H_n) + \underset{z=m}{\text{Res}}(-H_n) + \underset{z=\bar{m}}{\text{Res}}(-H_n)$$

$$\text{where } -H_n(z) = z / (z^n [z - b][z - m][z - \bar{m}]).$$

The sum of these three residues is

$$\begin{aligned} & \frac{b}{b^n [b - m][b - \bar{m}]} + \frac{m}{m^n [m - b][m - \bar{m}]} + \frac{\bar{m}}{\bar{m}^n [\bar{m} - b][\bar{m} - m]} \\ &= \mathbf{Q} \cdot p^n + \mathbf{U} \cdot w^n + \bar{U} \cdot \bar{w}^n \\ &= \mathbf{Q} \cdot p^n + \text{Re}(\mathbf{U} \cdot 2w^n), \quad \text{where} \end{aligned}$$

$$\begin{aligned} \mathbf{Q} &= \frac{1/p}{[1/p - 1/w][1/p - 1/\bar{w}]} = \frac{pw\bar{w}}{[w - p][\bar{w} - p]} \xrightarrow{\text{note}} \frac{1}{[p - w][p - \bar{w}]}, \\ \mathbf{U} &= \frac{1/w}{[1/w - 1/p][1/w - 1/\bar{w}]} = \frac{pw\bar{w}}{[p - w][\bar{w} - w]} \xrightarrow{\text{note}} \frac{1}{[p - w][\bar{w} - w]}. \end{aligned}$$

12b: **Theorem.** For each $n \in \mathbb{Z}$:

$$\mathbf{t}_n = \frac{p^n}{[p - w][\bar{w} - w]} + \text{Re}\left(\frac{2w^n}{[p - w][\bar{w} - w]}\right). \diamond$$

After typing up this TRIBONACCI task, now I find the following webpage:

https://en.wikipedia.org/wiki/Generalizations_of_Fibonacci_numbers#Tribonacci_numbers

Number theory of Fib sequence

T.fol problem and soln is from [Stack Exchange](#).

13: **Five-Fib thm.** For prime p , fibonacci number f_{p-1} is divisible by p IFF $p \equiv_5 \pm 1$. \diamond

Note. The result holds for $p=2$, as both parts fail. Henceforth, p is an odd prime. \square

Pf (\Leftarrow). Suppose $p \equiv_5 +1$ or $p \equiv_5 -1$. Then p is a 5-QR. Courtesy Quadratic Reciprocity, 5 is a p -QR; this, since 5 is 4POS.

Henceforth, \equiv means \equiv_p , and $\langle \cdot \rangle$ means $\langle \cdot \rangle_p$.

Let σ be such that $\sigma^2 \equiv 5$; use $\hat{\sigma} := \langle 1 \div \sigma \rangle$ for its reciprocal. Let $\hat{2} := \langle 1 \div 2 \rangle$. Binet's formula for f_n is

$$f_n = \frac{1}{\varphi} \cdot [\alpha^n - \beta^n], \quad \text{where} \\ \alpha, \beta = (1 \pm \varphi) \cdot \frac{1}{2}, \quad \text{where } \varphi := \sqrt{5}.$$

The Idea. Follow Binet by mimicking f_n in \mathbb{Z}_p :

$$* \begin{aligned} g_n &:= \hat{\sigma} \cdot \left[([1+\sigma] \cdot \hat{2})^n - ([1-\sigma] \cdot \hat{2})^n \right] \\ &= \hat{\sigma} \cdot \left[(1+\sigma)^n - (1-\sigma)^n \right] \cdot \hat{2}^n. \end{aligned}$$

Fermat's Little Thm. There is no odd prime p for which $5 \equiv 1$. Hence the above $\sigma \not\equiv \pm 1$. Thus $1+\sigma$ and $1-\sigma$ are p -units, as is $\hat{2}$.

We may thus apply FLiT to conclude that

$$**: \begin{aligned} g_{p-1} &\equiv \hat{\sigma} \cdot \left[([1+\sigma] \cdot \hat{2})^{p-1} - ([1-\sigma] \cdot \hat{2})^{p-1} \right] \\ &\stackrel{\text{FLiT}}{\equiv} \hat{\sigma} \cdot [1 - 1] = 0. \end{aligned}$$

Showing $g_N \equiv f_N$. ISTE Establish $2^N g_N \stackrel{?}{\equiv} 2^N f_N$, since $2 \perp p$. To this end, let H be the largest integer with $2H+1 \leq N$. Then $\varphi \cdot 2^N f_N$ equals

$$\begin{aligned} (1+\varphi)^N - (1-\varphi)^N &= \sum_{j=0}^N \binom{N}{j} \cdot [1 - (-1)^j] \varphi^j \\ &= \sum_{d=0}^H \binom{N}{2d+1} \cdot 2 \varphi^{2d+1} \\ &= 2\varphi \cdot \sum_{d=0}^H \binom{N}{2d+1} \cdot 5^d. \end{aligned}$$

Hence $2^N f_N$ equals

$$\frac{1}{\varphi} \cdot \left[(1+\varphi)^N - (1-\varphi)^N \right] = 2 \sum_{d=0}^H \binom{N}{2d+1} \cdot 5^d \stackrel{\text{note}}{\equiv} 2N.$$

The same algebra shows $2^N g_N$ is mod-5 congruent to

$$\hat{\sigma} \cdot \left[(1+\sigma)^N - (1-\sigma)^N \right] \equiv 2 \sum_{d=0}^H \binom{N}{2d+1} \cdot 5^d \equiv 2N. \diamond$$

Pf (\Rightarrow). $\textcolor{red}{?}$ \diamond

13a: **Corollary.** For n a natnum: $2^n \cdot f_n \equiv_5 2n$. \diamond