

# Fibonacci sequences

Jonathan L.F. King  
squash@ufl.edu  
Webpage <http://squash.1gainesville.com/>  
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See also [Problems/Algebra/LinearAlg/linear-recurr.tex](#)

**GCD.** We'll sometimes use a red under-bracket for *greatest common divisor*, e.g,  $\underline{n, m} := \text{GCD}(n, m)$ .

**Prolegomenon.** The famous *Fibonacci sequence*  $\vec{f} := (f_n)_{n=-\infty}^{\infty}$  is defined by  $f_0 := 0$ ,  $f_1 := 1$  and

$$1a: \quad f_{n+1} = f_n + f_{n-1},$$

producing this doubly- $\infty$  sequence:

$n \dots -5 \ -4 \ -3 \ -2 \ -1 \ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \dots$   
 $f_n \dots 5 \ -3 \ 2 \ -1 \ 1 \ 0 \ 1 \ 1 \ 2 \ 3 \ 5 \ 8 \ 13 \ 21 \ 34 \ 55 \dots$

Let  $\alpha$  and  $\beta$  be the positive and negative roots of the *characteristic polynomial* of  $\vec{f}$ , which is

$$\text{Fib}(x) := x^2 - x - 1 \stackrel{\text{note}}{=} [x - \alpha][x - \beta]. \text{ So}$$

$$1b: \quad \alpha + \beta = 1 \text{ and } \alpha \cdot \beta = -1. \text{ Moreover} \\ \alpha^2 = \alpha + 1, \beta^2 = \beta + 1 \text{ and } \alpha, \beta = \frac{1}{2}[1 \pm \sqrt{5}].$$

For future reference,

$$1c: \quad \alpha > 1 > |\beta|.$$

**Binet Formula.** Let  $\varphi := \sqrt{5}$ . [Mnemonic “phi” to evoke “five”.] Our  $\vec{f}$  is some linear combination  $A \cdot [n \mapsto \alpha^n] + B \cdot [n \mapsto \beta^n]$ . Easily,  $-B = A = \frac{1}{\varphi}$ , so

$$2: \quad \forall n \in \mathbb{Z}: \quad f_n = \frac{1}{\varphi} \cdot [\alpha^n - \beta^n],$$

since this formula gives correct values for  $f_0$  and  $f_1$ .

**2a: Lemma.** *Courtesy Binet:*

$$\dagger: \quad \forall n \in \mathbb{Z}: \quad \alpha^n = f_n \alpha + f_{n-1}. \text{ Also,}$$

$$\ddagger: \quad 5\alpha^2 = [\alpha + 2]^2.$$

*Each of these holds with  $\alpha$  replaced by  $\beta$ .*  $\diamond$

**Pf of  $(\dagger)$ .** We prove  $(\dagger)$  for  $n$  positive. The base case is  $\alpha^1 = 1 \cdot \alpha + 0 = f_1 \cdot \alpha + f_0$ . Inductively

$$\begin{aligned} \alpha^{n+1} &= \alpha \cdot [f_n \alpha + f_{n-1}] = f_n \alpha^2 + f_{n-1} \alpha \\ &= f_n [\alpha + 1] + f_{n-1} \alpha \\ &= [f_n + f_{n-1}] \alpha + f_n \end{aligned}$$

which equals  $f_{n+1} \alpha + f_n$ , as desired.  $\blacklozenge$

**Pf  $(\ddagger)$ .** Squaring,  $[\alpha + 2]^2 = \alpha^2 + 4[\alpha + 1] = \alpha^2 + 4\alpha^2$ .  $\blacklozenge$

### Recurrence doubling

A fibonacci-like sequence  $\vec{z} := (z_n)_{n=0}^\infty$  is defined via (possibly complex) numbers  $\mathcal{P}$  and  $\mathcal{S}$ , by

$$3.1: \quad z_{n+2} := \mathcal{S}z_{n+1} - \mathcal{P}z_n,$$

and some initial condition  $(z_1, z_0)$ . With

$$3.2: \quad \mathbf{G} := \begin{bmatrix} \mathcal{S} & -\mathcal{P} \\ 1 & 0 \end{bmatrix}, \quad \text{then } \begin{bmatrix} z_{n+1} \\ z_n \end{bmatrix} = \mathbf{G}^n \cdot \begin{bmatrix} z_1 \\ z_0 \end{bmatrix},$$

for each integer  $n$ .

A number  $r \neq 0$  engenders seq  $n \mapsto r^n$ . It satisfies recurrence (3.1) IFF  $r$  is a root of polynomial

$$f(x) := x^2 - \mathcal{S}x + \mathcal{P} \stackrel{\text{factored}}{=} [x - \mu][x - \nu].$$

Equating coeffs in the polynomial gives:

$$3.3: \quad \mu + \nu = \mathcal{S}. \quad (\text{Sum})$$

$$\mu \cdot \nu = \mathcal{P}. \quad (\text{Product})$$

$$3.4: \quad \mu^2 = \mathcal{S}\mu - \mathcal{P} \quad \text{and} \quad \nu^2 = \mathcal{S}\nu - \mathcal{P}.$$

Henceforth, we require  $\mu \neq \nu$ , i.e.  $\mathcal{S}^2 \neq 4\mathcal{P}$ ; this, since  $\text{Discr}(f) \stackrel{\text{note}}{=} \mathcal{S}^2 - 4\mathcal{P}$ . We want  $\mu \neq \nu$  so that every seq satisfying (3.1) has its  $n^{\text{th}}$ -term equal a linear combination of  $\mu^n$  and  $\nu^n$ .

**Doubling.** Sequences  $(z_e)_{e \text{ even}}$  and  $(z_d)_{d \text{ odd}}$  will satisfy some two-term linear recurrence. [The same recurrence.] We seek numbers  $\hat{\mathcal{S}}$  and  $\hat{\mathcal{P}}$  such that for  $r$  equaling either  $\mu^2$  or  $\nu^2$ ,

$$r^2 = \hat{\mathcal{S}}r^1 - \hat{\mathcal{P}}r^0.$$

I.e, that polynomial

$$x^2 - \hat{\mathcal{S}}x + \hat{\mathcal{P}} \stackrel{\text{factors}}{=} [x - \mu^2] \cdot [x - \nu^2].$$

Squaring eqn  $\mu^2 = \mathcal{S}\mu - \mathcal{P}$  gives

$$\mu^4 = \mathcal{S}^2\mu^2 - 2\mathcal{P}\mathcal{S}\mu + \mathcal{P}^2.$$

As  $\mathcal{S}\mu = \mu^2 + \mathcal{P}$ , so  $2\mathcal{P}\mathcal{S}\mu = 2\mathcal{P}\mu^2 + 2\mathcal{P}^2$ . Thus

$$\mu^4 = \underbrace{[\mathcal{S}^2 - 2\mathcal{P}]}_{\hat{\mathcal{S}}} \mu^2 - \underbrace{\mathcal{P}^2}_{\hat{\mathcal{P}}}.$$

**3.5: Doubling thm.** Sequence  $\vec{z} := (z_n)_{n=0}^\infty$  satisfies recurrence

$$z_{n+2} := \mathcal{S}z_{n+1} - \mathcal{P}z_n, \quad \text{where } \mathcal{S}, \mathcal{P} \in \mathbb{C}.$$

Then sequence  $(z_{2n})_{n \in \mathbb{N}}$  and  $(z_{2n+1})_{n \in \mathbb{N}}$  each satisfy recurrence

$$\square_{n+2} = \hat{\mathcal{S}} \cdot \square_{n+2} - \hat{\mathcal{P}} \cdot \square_n$$

where  $\hat{\mathcal{S}} := \mathcal{S}^2 - 2\mathcal{P}$  and  $\hat{\mathcal{P}} := \mathcal{P}^2$ .

[If  $\mathcal{P} \neq 0$ , then  $\vec{z}$  can be extended backwards to  $\vec{z} := (z_n)_{n=-\infty}^\infty$ , and the doubling result holds.]  $\diamond$

**Proof.** The foregoing argument used  $\mathcal{S}^2 \neq 4\mathcal{P}$ . Since sequence-values vary continuously with  $\mathcal{S}$  and  $\mathcal{P}$ , it suffices to obtain an  $\mathcal{S}^2 = 4\mathcal{P}$  pair as a limit of pairs where the inequality holds.  $\blacklozenge$

**Fib example.** The Fib-seq has  $\mathcal{S} = 1$  and  $\mathcal{P} = -1$ . Hence  $\hat{\mathcal{S}} = 1^2 - 2(-1) = 3$ , and  $\hat{\mathcal{P}} = [-1]^2$ . So  $(f_{2n})_{n \in \mathbb{Z}}$  and  $(f_{2n+1})_{n \in \mathbb{Z}}$  each satisfy  $z_{n+2} = 3z_{n+1} - z_n$ .

Starting at any index, then taking every 4<sup>th</sup>-term, gives a seq satisfying  $z_{n+2} = 7z_{n+1} - z_n$ , since  $3^2 - 2 \cdot 1 = 7$  and  $1^2 = 1$ .  $\square$

**3.6: Speed-up thm.** Sequence  $\vec{z} := (z_n)_{n=0}^\infty$  satisfies recurrence  $z_{n+2} := \mathcal{S}z_{n+1} - \mathcal{P}z_n$ .

For posint  $K$  and natnum  $\mathbf{b}$ , subsequence  $(z_{nK+\mathbf{b}})_{n=0}^\infty$  satisfies  $\square_{n+2} = \hat{\mathcal{S}} \cdot \square_{n+2} - \hat{\mathcal{P}} \cdot \square_n$ , where  $x^2 - \hat{\mathcal{S}}x + \hat{\mathcal{P}}$  is the characteristic polynomial of  $K^{\text{th}}$ -power

$$\mathbf{G}^K = \begin{bmatrix} \mathcal{S} & -\mathcal{P} \\ 1 & 0 \end{bmatrix}^K. \quad \diamond$$

**Pf.** Use  $x^2 - \hat{\mathcal{S}}x + \hat{\mathcal{P}}$  for the charpoly of  $\mathbf{M} := \mathbf{G}^K$ . Cayley-Hamilton asserts  $\mathbf{M}^2 = \hat{\mathcal{S}} \cdot \mathbf{M} - \hat{\mathcal{P}} \cdot \mathbf{I}$ . Applying both sides to  $\text{colvec} \begin{bmatrix} z_1 \\ z_0 \end{bmatrix}$ , the bottom entry asserts

$$z_{2K} = \hat{\mathcal{S}} \cdot z_K - \hat{\mathcal{P}} \cdot z_0. \quad \blacklozenge$$

**4: Theorem.** For each integer  $N$ :

$$4a: \quad [f_N]^2 + [f_{N-1}]^2 = f_{2N-1}. \quad \diamond$$

**Inefficient proof.** Always,  $\text{LhS}(4a)$  is non-negative. And  $\text{RhS}(4a)$  is non-neg, even when  $N \in \mathbb{Z}_-$ , since  $f_{\text{Odd}}$  is always non-neg. So ISTProve that the squares of  $\text{LhS}(4a)$  and  $\text{RhS}(4a)$  are equal. To this end, define

$$4b: \quad \begin{aligned} \mathcal{L} &:= \varphi^2 \cdot ([f_N]^2 + [f_{N-1}]^2) \quad \text{and} \\ \mathcal{R} &:= \varphi^2 \cdot [f_{2N-1}] \stackrel{\text{note}}{=} \varphi \cdot [\alpha^{2N-1} - \beta^{2N-1}]. \end{aligned}$$

**Leftside.** By (2),  $\varphi \cdot f_N = \alpha^N - \beta^N$ . So  $\varphi^2 \cdot f_N^2$  equals

$$\alpha^{2N} + \beta^{2N} - 2[\alpha\beta]^N.$$

But  $\alpha\beta = -1$ , and  $N$  and  $N-1$  have opposite parities. Thus  $\mathcal{L}$  equals

$$\begin{aligned} \alpha^{2N} + \beta^{2N} + \alpha^{2[N-1]} + \beta^{2[N-1]} \\ = \alpha^{2[N-1]}[\alpha^2 + 1] + \dots, \end{aligned}$$

where the “ $\dots$ ” represents a copy of all the  $\alpha$ -terms to its left, but with “ $\alpha$ ” replaced by “ $\beta$ ”.

By (1b), note,  $\alpha^2 + 1 = \alpha + 2$ . Thus

$$\mathcal{L} = \alpha^{2[N-1]}[\alpha + 2] + \beta^{2[N-1]}[\beta + 2].$$

Squaring  $\mathcal{L}$  will give twice this cross-term:

$$\begin{aligned} [\alpha\beta]^{2[N-1]} \cdot [\alpha + 2][\beta + 2] &= 1 \cdot [\alpha + 2][\beta + 2] \\ &= \alpha\beta + 2[\alpha + \beta] + 4 \\ &= -1 + 2 + 4 = 5. \end{aligned}$$

Also note  $[\alpha + 2]^2 = \alpha^2 + 4[\alpha + 1] = 5\alpha^2$ . Thus

$$\mathcal{L}^2 = \alpha^{4[N-1]} \cdot 5\alpha^2 + \dots + 5 \cdot 2.$$

Consequently

$$4c: \quad \frac{1}{5} \cdot \mathcal{L}^2 = \alpha^{4N-2} + \beta^{4N-2} + 2.$$

**Rightside of (4a).** Square  $\mathcal{R}$  and divide by 5. Since  $5 = \varphi^2$ ,

$$\frac{1}{5} \cdot \mathcal{R}^2 = [\alpha^{2N-1} - \beta^{2N-1}]^2.$$

The cross-term is  $-2[\alpha\beta]^{2N-1} = -2 \cdot [-1]^{2N-1} = 2$ , since  $2N-1$  is odd. We have thus shown that

$$4d: \quad \frac{1}{5} \cdot \mathcal{R}^2 = \alpha^{4N-2} + \beta^{4N-2} + 2.$$

And this equals  $\text{RhS}(4c)$ , as desired.  $\diamond$

*Ahem.* It’s a proof, but the next is a *prettier proof*.  $\square$

### Dot-product proof

Fix three integers  $S = j + k$ . Evidently the dot-product

$$\begin{aligned} f_{j+1}f_k + f_jf_{k-1} &= f_{j+1}[f_{k-1} + f_{k-2}] + f_jf_{k-1} \\ &= [f_{j+1} + f_j]f_{k-1} + f_{j+1}f_{k-2} \\ &= f_{j+2}f_{k-1} + f_{j+1}f_{k-2}. \end{aligned}$$

This last dot-product is the same as the first, but with “ $j + k = S$ ” replaced by “ $[j+1] + [k-1] = S$ ”. Thus, for all  $j \in \mathbb{Z}$ , expression  $f_{j+1} \cdot f_k + f_j \cdot f_{k-1}$  depends only on  $S$ .

**5: Theorem.** For all triples of integers  $S = j + k$ :

$$5a: \quad f_{j+1} \cdot f_k + f_j \cdot f_{k-1} = f_S.$$

In particular, (4a) holds.  $\diamond$

**Proof.** Setting  $j := 0$  in  $\text{LhS}(5a)$  results in

$$f_1f_S + f_0f_{S-1} = 1 \cdot f_S + 0 \cdot f_{S-1} = f_S.$$

Hence (5a). To obtain (4a) from (5a), set  $S := 2N-1$  and  $j := N-1$ .  $\diamond$

**Alt Pf.** Looking ahead to matrix  $A := \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$  from Cassini, the  $(1, 2)$ -entry of matrix  $A^S = A^j A^k$  is  $(*)$ .  $\diamond$

**5b: Duplication identity.** For each integer  $n$ ,

$$*: \quad f_{2n} = f_n \cdot [f_{n+1} + f_{n-1}] = f_n \cdot [f_n + 2f_{n-1}]. \quad \diamond$$

**Proof.** Apply (5a) with  $j := n-1$  and  $k := n+1$ . Then  $f_{2n} = f_n f_{n+1} + f_{n-1} f_n = f_n \cdot [f_{n+1} + f_{n-1}]$ . So  $f_{2n}$  equals

$$f_n \cdot [f_n + f_{n-1}] + f_{n-1} = \text{RhS}(*) . \quad \diamond$$

**5c: Cassini's identity.** For all  $n$ ,

$$\mathcal{L}: \quad f_{n+1} \cdot f_{n-1} = f_n^2 + [-1]^n. \quad \diamond$$

**Pf.** Let  $\mathbf{v}_n := \begin{bmatrix} f_n \\ f_{n-1} \end{bmatrix}$ . With  $\mathbf{A} := \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$ , our recurrence says  $\mathbf{A}\mathbf{v}_n = \mathbf{v}_{n+1}$ . Id-matrix  $\mathbf{I} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} f_1 & f_0 \\ f_0 & f_{-1} \end{bmatrix}$ ; i.e  $\mathbf{A}^0$  is  $2 \times 2$  matrix  $\begin{bmatrix} \mathbf{v}_1 & \mathbf{v}_0 \end{bmatrix}$ . Thus  $\mathbf{A}^n$  equals  $\begin{bmatrix} f_{n+1} & f_n \\ f_n & f_{n-1} \end{bmatrix}$ . Equation  $\text{Det}(\mathbf{A}^n) = [\text{Det}(\mathbf{A})]^n$  yields  $(\mathcal{L})$ .  $\blacklozenge$

**Periodicity mod a prime.** Working mod-5, note

$$(f_{10}, f_{11}) = (55, 89) \equiv_5 (0, -1) = [-1] \cdot (f_0, f_1).$$

Mod-5, then, Fib-seq  $\vec{f}$  has a nega-period of length 10, whence  $\vec{f}_{\text{Mod-5}}$  has 20 as a period:  $f_{20+n} \equiv_5 f_n$ .

[Eventually adjoin material wrt other primes.]

**5d: 1-3-1 Fib lemma.** Fibonacci sequence  $\vec{f}$  satisfies

$$\mathcal{L}\mathcal{L}: \quad [f_{n+1}]^2 - 3[f_n]^2 + [f_{n-1}]^2 = 2 \cdot [-1]^n. \quad \diamond$$

**Proof.** Squaring  $f_n = f_{n+1} - f_{n-1}$  implies

$$\begin{aligned} f_n^2 &= f_{n+1}^2 + f_{n-1}^2 - 2f_{n+1} \cdot f_{n-1} \\ &\stackrel{\text{Cassini}}{=} f_{n+1}^2 + f_{n-1}^2 - 2[f_n^2 + [-1]^n] \\ &= f_{n+1}^2 + f_{n-1}^2 - 2f_n^2 - 2[-1]^n. \quad \blacklozenge \end{aligned}$$

**6: Lemma.**  $\forall n \in \mathbb{N}: [f_1]^2 + [f_2]^2 + \dots + [f_n]^2 = f_n f_{n+1}. \quad \diamond$

**Pf.** Easy induction. Even nicer,  $f_n f_{n+1}$  is the area of a  $f_n \times f_{n+1}$  business card. Decompose it into squares.  $\blacklozenge$

## Divisibility

Integer-sequence  $\vec{\sigma} = (\sigma_1, \sigma_2, \dots)$  is a **divisibility sequence** [*div-seq*] if  $\forall_{\text{indices } j, k \in \mathbb{Z}_+}$ : If  $j \blacktriangleright k$  then  $\sigma_j \blacktriangleright \sigma_k$ .

Our  $\vec{\sigma}$  is a **strong div-seq** if  $\forall j, k$ :

$$\dagger: \quad \sigma_{\text{GCD}(j, k)} = \text{GCD}(\sigma_j, \sigma_k).$$

Why does *strong div-seq* implies *div-seq*? When  $j \blacktriangleright k$ , then  $\text{GCD}(j, k) = j$ . Hence  $\text{GCD}(\sigma_j, \sigma_k) = \sigma_j$ . And this implies  $\sigma_j \blacktriangleright \sigma_k$ .

Strong div-seq implies, for each list  $k_1, \dots, k_N$ ,

$$\dagger: \quad \sigma_{\text{GCD}(k_1, \dots, k_N)} = \text{GCD}(\sigma_{k_1}, \dots, \sigma_{k_N}).$$

Using underline for GCD, note  $a, c, e$  =  $a, c$ ,  $e$ . Hence,

$$\begin{aligned} \sigma_{\underline{a, c, e}} &= \sigma_{\underline{a, c}, e} \xrightarrow{\text{by } (\dagger)} \sigma_{\underline{a, c}, \sigma_e} \\ &\xrightarrow{\text{by } (\dagger)} \sigma_{\underline{\sigma_a, \sigma_c}, \sigma_e} \\ &= \sigma_{\underline{\sigma_a, \sigma_c, \sigma_e}}. \end{aligned}$$

**8.1: Lemma.** For integers  $b > a \geq 0$ , values  $\sigma_n := b^n - a^n$  form a divisibility sequence.  $\diamond$

**Proof.** Given  $j \blacktriangleright k$ , write  $k = \rho \cdot j$ , and set  $B := b^j$  and  $A := a^j$ . So

$$\frac{\sigma_k}{\sigma_j} = \frac{B^\rho - A^\rho}{B - A} = \sum_{\substack{v+u=\rho-1 \\ v, u \geq 0}} B^v A^u. \quad \diamond$$

**8.2: Obs.** Div-seq  $\sigma_n = b^n - a^n$  need not be a strong div-sequence. E.g, with  $b := 4$  and  $a := 2$ , note

$$\text{GCD}(\sigma_6, \sigma_4) = \text{GCD}(4032, 240) = 48.$$

Yet  $\sigma_{\text{GCD}(6, 4)} = \sigma_2 = 12$ .  $\square$

**8.3: ? Fact?** For integers  $b > a \geq 1$ , if  $b \perp a$  then sequence  $\sigma_n := b^n - a^n$  is a strong divisibility-sequence.  $\diamond$

**Proof. ??** See proof of (11).  $\diamond$

**9: Theorem.** Fibonacci seq  $\vec{f} = (f_1, f_2, f_3, \dots)$  is a divisibility-sequence.  $\diamond$

**Pf.** Fix posint  $P$ . For  $n = 1, 2, \dots$ , we prove  $f_P \mid f_{nP}$ . The base-case is  $f_P \mid f_{1 \cdot P}$ .

We apply Thm 5,

$$*: \quad f_{j+k} = f_{j+1} \cdot f_k + f_j \cdot f_{k-1},$$

with  $j := nP$  and  $k := P$ . Inductively, quotient  $Q := f_{nP}/f_P$  is an integer. Our (\*) says

$$\begin{aligned} f_{[n+1]P} &= f_{nP+1}f_P + f_{nP}f_{P-1} \\ &= f_{nP+1}f_P + Qf_Pf_{P-1} \\ &= f_P \cdot [f_{nP+1} + Qf_{P-1}]. \end{aligned}$$

And  $f_{nP+1} + Qf_{P-1}$  is an integer.  $\diamond$

**10.1: Prop.** For integers  $\delta, \beta, D$ : If  $\delta \mid D$ , then

$$\text{GCD}(\delta + \beta, D) = \text{GCD}(\beta, D). \quad [\text{Exercise}] \quad \diamond$$

**10.2: Prop.** For  $n$  an integer:  $f_{n+1} \perp f_n$ . [Exercise]  $\diamond$

**11: Fib strong-div.** Divide  $u \neq 0$  into  $v$  [both integers] to get quotient and remainder,  $v = qu + r$ . Then

$$11a: \quad \text{GCD}(f_v, f_u) = \text{GCD}(f_u, f_{v-u});$$

$$11b: \quad \text{GCD}(f_v, f_u) = \text{GCD}(f_u, f_r);$$

$$11c: \quad \text{GCD}(f_v, f_u) = f_{\text{GCD}(v,u)}.$$

Thus  $f_1, f_2, f_3, \dots$  is a strong divisibility-sequence.  $\diamond$

**Pf 11a.** We apply Thm 5 with  $j := v - u$  and  $k := u$ . So

$$f_v = f_{v-u+1}f_u + f_{v-u}f_{u-1}.$$

Since  $f_u$  divides  $f_{v-u+1}f_u$ , our Prop 10.1 says

$$\underline{f_v, f_u} = \underline{f_{v-u}f_{u-1}, f_u} = \underline{f_{v-u}, f_u},$$

since  $f_{u-1} \perp f_u$ . [Argument works also for  $u=0$ .]  $\diamond$

**Pf 11b.** Applying (11a)  $q$  times gives  $\underline{f_v, f_u} = \underline{f_u, f_{v-qu}}$   $\diamond$

**Pf 11c.** Recall the update rule in the Euclidean algorithm (Lightning Bolt) when seeding LBolt with  $r_0 := v$  and  $r_1 := u$ . Observe that the  $r$  of decomposition  $v = qu + r$  is the  $r_2$  of LBolt. Thus (11b) says

$$\underline{f_{r_0}, f_{r_1}} = \underline{f_{r_1}, f_{r_2}}.$$

But *this* is the update rule when seeded with  $f_{r_0}$  and  $f_{r_1}$ . Consequently, letting  $\underline{g} := \underline{v, u}$ ,

$$\underline{f_v, f_u} = \underline{f_g, f_0} = \underline{f_g, 0} \stackrel{\text{note}}{=} f_g. \quad \diamond$$

## Can complex analysis Fib?

What is the RoC of

$$\dagger: \quad \mathcal{F}(z) := \sum_{n \geq 0} f_n z^n,$$

the *generating function* of the fibonacci sequence?

Recurrence  $f_{n+2} = f_{n+1} + f_n$ , and that  $f_n \geq 0$  for  $n \geq 0$  show that sequence  $\vec{f}$  is non-decreasing. For  $n \geq 2$ , the sequence is positive, so recurrence (1a) gives

$$\begin{aligned} 2f_{n-2} &\leq f_n \leq 2f_{n-1}. \quad \text{Hence} \\ \text{Const} \cdot [\sqrt{2}]^n &\leq f_n \leq \text{Const} \cdot 2^n. \quad \text{Thus} \\ \frac{1}{\sqrt{2}} &\geq \text{RoC}(\mathcal{F}) \geq \frac{1}{2}. \end{aligned}$$

In particular,  $\mathcal{F}$  is analytic in a nbhd of zero.

**Shh...** In fact we know that

$$\begin{aligned} ** : \quad f_n &= \frac{1}{\sqrt{5}} \cdot [\alpha^n - \beta^n], \quad \text{where} \\ \alpha &= \frac{1}{2}[1 + \sqrt{5}] \quad \text{and} \quad \beta = \frac{1}{2}[1 - \sqrt{5}]. \end{aligned}$$

Here,  $\alpha$  and  $\beta$  are the *golden/silver* ratios, respectively. We thus have that

$$\text{Shh!}: \quad \text{RoC}(\mathcal{F}) = 1/\alpha = \frac{2}{1 + \sqrt{5}}.$$

**Source.** Online book **A FIRST COURSE IN COMPLEX ANALYSIS**, made freely available by Matthias Beck, Gerald Marchesi, Dennis Pixton & Lucas Sbalka, has the below problem in §10.3 of their text.  $\square$

**A formula for  $\mathcal{F}$ .** [Below, I'll use  $\mathcal{F}$  to abbreviate  $\mathcal{F}(z)$ .]

By defn,

$$\begin{aligned} z^0 \cdot \sum_{n \geq 0} f_{n+2} \cdot z^{n+2} &= \mathcal{F} - f_0 \cdot z^0 - f_1 \cdot z^1 \stackrel{\text{note}}{=} \mathcal{F} - z; \\ z^1 \cdot \sum_{n \geq 0} f_{n+1} \cdot z^{n+1} &= z \cdot [\mathcal{F} - f_0 \cdot z^0] \stackrel{\text{note}}{=} z \cdot \mathcal{F}; \\ z^2 \cdot \sum_{n \geq 0} f_n \cdot z^n &= z^2 \cdot \mathcal{F}. \end{aligned}$$

From the top equality, we subtract the other two. The lefthand side of the result is

$$\sum_{n \geq 0} [f_{n+2} - f_{n+1} - f_n] z^{n+2} \stackrel{\text{note}}{=} 0,$$

since  $f_{n+2} = f_{n+1} + f_n$ . From the righthand sides, then,  $0 = [1 - z - z^2]\mathcal{F} - z$ . Hence

$$\dagger: \quad \mathcal{F}(z) = \frac{z}{D(z)}, \quad \text{where} \quad D(z) := 1 - z - z^2.$$

Easily,  $D(z) = z^2 \cdot \text{Fib}(\frac{1}{z}) = z^2 \cdot [\frac{1}{z} - \alpha][\frac{1}{z} - \beta]$ . Rewriting,  $D(z) = [\alpha z - 1][\beta z - 1]$ . Thus

$$D(z) = \alpha\beta[z - \frac{1}{\alpha}][z - \frac{1}{\beta}] \stackrel{\text{note}}{=} -1 \cdot [z + \beta][z + \alpha].$$

**Residues.** With  $H_n(z) := \frac{1}{z^n D(z)} \stackrel{\text{def}}{=} \frac{\mathcal{F}(z)}{z^{n+1}}$ , note

$$\begin{aligned} \text{Res}_{z=0}(H_n) &= \text{Res}_{z=0}\left(\frac{\mathcal{F}(z)}{z^{n+1}}\right) \\ &\stackrel{\text{PlexNotes}}{\stackrel{P.30}{=}} \frac{1}{n!} \cdot \mathcal{F}^{(n)}(0) = f_n. \end{aligned}$$

Looking ahead,  $\frac{1}{\beta^n} - \frac{1}{\alpha^n} = \frac{\alpha^n - \beta^n}{[\alpha\beta]^n} = \frac{\alpha^n - \beta^n}{[-1]^n}$ . So

$$\P: \quad \frac{1}{[-\beta]^n} - \frac{1}{[-\alpha]^n} = \alpha^n - \beta^n.$$

**Contour integral.** Let  $\mathbf{C}_r := \text{Sph}_r(0)$ . Since  $n \geq 0$ , the degree of  $z^n D(z)$  is at least 2, whence

$$\lim_{r \nearrow \infty} \frac{1}{2\pi i} \oint_{\mathbf{C}_r} H_n = 0.$$

As the singularities of  $H_n$  are  $0, -\beta, -\alpha$ ,

$$\begin{aligned} 0 &= \text{Res}(H_n, 0) + \text{Res}(H_n, -\beta) + \text{Res}(H_n, -\alpha), \quad \text{i.e} \\ * : \quad f_n &= \text{Res}(-H_n, -\beta) + \text{Res}(-H_n, -\alpha). \end{aligned}$$

Writing  $-H_n(z) = \frac{1/[z^n(\alpha + z)]}{z - \beta}$ , the CIF gives

$$\begin{aligned} \text{Res}(-H_n, -\beta) &= 1/[-\beta]^n[\alpha - \beta]. \quad \text{Similarly,} \\ \text{Res}(-H_n, -\alpha) &= 1/[-\alpha]^n[\beta - \alpha]. \end{aligned}$$

Their sum equals

$$\frac{1}{\alpha - \beta} \cdot \left[ \frac{1}{[-\beta]^n} - \frac{1}{[-\alpha]^n} \right] \stackrel{\text{by } (\P)}{=} \frac{1}{\alpha - \beta} \cdot [\alpha^n - \beta^n].$$

Plugging this into (\*), yields (\*\*) –which we supposedly didn't know.

## Tribonacci sequence

The bi-infinite **Trib sequence** is

$$\begin{aligned} * : \quad & (t_0, t_1, t_2) := (0, 0, 1) \quad \text{and} \\ & t_{n+3} := t_{n+2} + t_{n+1} + t_n, \end{aligned}$$

for  $n \in \mathbb{Z}$ . The resulting Trib sequence is

$$\begin{aligned} n: & \dots -9 -8 -7 -6 -5 -4 -3 -2 -1 0 1 2 3 4 5 6 7 \dots \\ t_n: & \dots 81 44 24 13 7 4 2 1 1 0 0 1 1 2 4 7 13 \dots \end{aligned}$$

A nz-complex  $\omega$  engenders seq  $n \mapsto \omega^n$ , which satisfies the recurrence part of (\*) IFF  $\omega$  is a zero of

$$\text{Trb}(z) := z^3 - z^2 - z - 1.$$

**Zeros of Trb().** As  $\text{Trb}(1) < 0 < \text{Trb}(2)$ , value  $\text{Trb}(p)$  is zero for some posreal  $1 < p < 2$ . The remaining Trb-zeros form a complex-conjugate pair  $w \neq \bar{w}$ .

To see this last, ISTShow that  $\text{Trb}|_{\mathbb{R}}$  has but one zero. Its derivative equals

$$\text{Trb}'(x) = 3x^2 - 2x - 1 = 3[z - 1][z + \tfrac{1}{3}].$$

At the  $\text{Trb}()$ -critical-points, each of  $\text{Trb}(1) = -2$  and  $\text{Trb}(\frac{-1}{3}) = \frac{-22}{27}$  is negative. Hence  $\text{Trb}(x)$  is negative for  $x < p$ , and is positive for  $x > p$ .

**12: Prop'n.** Sum  $w + \bar{w} + p = 1$ , product  $w \cdot \bar{w} \cdot p = 1$ . Also,  $p > 1 > |w| = |\bar{w}| = \frac{1}{\sqrt{p}}$ .  $\diamond$

**Pf.** Write  $z^3 - 1z^2 - z - 1 = [z - w][z - \bar{w}][z - p]$  and equate coeffs.  $\blacklozenge$

**Trib GF.** The generating function of Trib-seq is

$$\dagger: \quad \mathcal{T}(z) := \sum_{n \geq 0} t_n z^n. \quad \square$$

**12a: Lemma.** The RoC( $\mathcal{T}$ ) is  $1/|p|$ . **Proof.** Exercise.  $\diamond$

*“Cheating”.* Cardano’s formula gives

$$\begin{aligned} p &= \tfrac{1}{3}[1 + S^+ + S^-], \\ w &= \tfrac{1}{6}[2 + \mu S^+ + \bar{\mu} S^-], \\ \bar{w} &= \tfrac{1}{6}[2 + \bar{\mu} S^+ + \mu S^-], \end{aligned}$$

where

$$\begin{aligned} \mu &:= -1 + i\sqrt{3}, & S^+ &:= \sqrt[3]{19 + 3\sqrt{33}}, \\ \bar{\mu} &:= -1 - i\sqrt{3}, & S^- &:= \sqrt[3]{19 - 3\sqrt{33}}. \end{aligned}$$

Values  $S^+ > S^-$  are real. And  $\text{Im}(w) > 0 > \text{Im}(\bar{w})$ . Approximately  $p \approx 1.839$  and  $w \approx -0.419 + 0.606i$ . However, we eschew this information FtTBeing.  $\square$

**A formula for  $\mathcal{T}$ .** [Below,  $\mathcal{T}$  abbreviates  $\mathcal{T}(z)$ .] By defn,

$$\begin{aligned} z^0 \cdot \sum_{n \geq 0} t_{n+3} \cdot z^{n+3} &= \mathcal{T} - t_0 \cdot z^0 - t_1 \cdot z^1 - t_2 \cdot z^2 \stackrel{\text{note}}{=} \mathcal{T} - z^2; \\ z^1 \cdot \sum_{n \geq 0} t_{n+2} \cdot z^{n+2} &= z \cdot [\mathcal{T} - t_0 \cdot z^0 - t_1 \cdot z^1] \stackrel{\text{note}}{=} z \cdot \mathcal{T}; \\ z^2 \cdot \sum_{n \geq 0} t_{n+1} \cdot z^{n+1} &= z^2 \cdot [\mathcal{T} - t_0 \cdot z^0] \stackrel{\text{note}}{=} z^2 \cdot \mathcal{T}; \\ z^3 \cdot \sum_{n \geq 0} t_n \cdot z^n &= z^3 \cdot \mathcal{T}. \end{aligned}$$

From the top equality, subtracting the lower three gives this LhS,

$$\sum_{n \geq 0} [t_{n+3} - t_{n+2} - t_{n+1} - t_n] z^{n+3} \stackrel{\text{note}}{=} 0,$$

by (\*). Hence  $0 = [1 - z - z^2 - z^3] \mathcal{T} - z^2$ , from the RhS. Consequently,

$$\dagger: \quad \mathcal{T}(z) = \frac{z^2}{E(z)}, \text{ where } E(z) := 1 - z - z^2 - z^3.$$

**Zeros of  $E()$ .** As  $E(z) = \text{Trb}(\frac{1}{z}) \cdot z^3$ , the zeros of  $E()$  are reciprocals  $b := 1/p$ ,  $m := 1/w$  and  $\bar{m}$ . Thus

$$E(z) = -1 \cdot [z - b][z - m][z - \bar{m}].$$



**Residues.** As  $\text{RoC}(\mathcal{T})$  is positive,  $\text{Res}_{z=0}(\frac{\mathcal{T}(z)}{z^{n+1}}) = \mathbf{t}_n$  for  $n=0, 1, \dots$ , courtesy  $(\dagger)$ .

Ratio  $\frac{\mathcal{T}(z)}{z^{n+1}} \xrightarrow{\text{by } (\dagger)} \frac{z}{z^n E(z)} =: H_n(z)$  is a rational function satisfying  $\text{Deg}(\text{Denominator}) \geq \text{Deg}(\text{Numerator}) + 2$ . Consequently,  $\lim_{r \nearrow \infty} \int_{\mathbf{C}_r} H_n(z) dz$  is zero.

For  $r$  large, circle  $\mathbf{C}_r$  encloses  $0, \mathbf{b}, \mathbf{m}, \overline{\mathbf{m}}$ , the four singularities of  $H_n$ . Hence their  $H_n$ -residues sum to zero. Recall that the  $z=0$  residue is  $\mathbf{t}_n$ . So...

$$\mathbf{t}_n = \text{Res}_{z=\mathbf{b}}(-H_n) + \text{Res}_{z=\mathbf{m}}(-H_n) + \text{Res}_{z=\overline{\mathbf{m}}}(-H_n)$$

where  $-H_n(z) = z / (z^n [z - \mathbf{b}][z - \mathbf{m}][z - \overline{\mathbf{m}}])$ .

The sum of these three residues is

$$\begin{aligned} & \frac{\mathbf{b}^n}{\mathbf{b}^n [\mathbf{b} - \mathbf{m}][\mathbf{b} - \overline{\mathbf{m}}]} + \frac{\mathbf{m}^n}{\mathbf{m}^n [\mathbf{m} - \mathbf{b}][\mathbf{m} - \overline{\mathbf{m}}]} + \frac{\overline{\mathbf{m}}^n}{\overline{\mathbf{m}}^n [\overline{\mathbf{m}} - \mathbf{b}][\overline{\mathbf{m}} - \mathbf{m}]} \\ &= \mathbf{Q} \cdot \mathbf{p}^n + \mathbf{U} \cdot \mathbf{w}^n + \overline{\mathbf{U}} \cdot \overline{\mathbf{w}}^n \\ &= \mathbf{Q} \cdot \mathbf{p}^n + \text{Re}(\mathbf{U} \cdot 2 \mathbf{w}^n), \quad \text{where} \\ \mathbf{Q} &= \frac{1/\mathbf{p}}{[1/\mathbf{p}-1/\mathbf{w}][1/\mathbf{p}-1/\overline{\mathbf{w}}]} = \frac{\mathbf{p}\mathbf{w}\overline{\mathbf{w}}}{[\mathbf{w}-\mathbf{p}][\overline{\mathbf{w}}-\mathbf{p}]} \stackrel{\text{note}}{=} \frac{1}{[\mathbf{p}-\mathbf{w}][\overline{\mathbf{p}}-\overline{\mathbf{w}}]}, \\ \mathbf{U} &= \frac{1/\mathbf{w}}{[1/\mathbf{w}-1/\mathbf{p}][1/\mathbf{w}-1/\overline{\mathbf{w}}]} = \frac{\mathbf{p}\mathbf{w}\overline{\mathbf{w}}}{[\mathbf{p}-\mathbf{w}][\overline{\mathbf{w}}-\mathbf{w}]} \stackrel{\text{note}}{=} \frac{1}{[\mathbf{p}-\mathbf{w}][\overline{\mathbf{w}}-\mathbf{w}]} \end{aligned}$$

**12b: Theorem.** For each  $n \in \mathbb{Z}$ :

$$\mathbf{t}_n = \frac{\mathbf{p}^n}{[\mathbf{p} - \mathbf{w}][\overline{\mathbf{w}} - \mathbf{w}]} + \text{Re}\left(\frac{2 \mathbf{w}^n}{[\mathbf{p} - \mathbf{w}][\overline{\mathbf{w}} - \mathbf{w}]}\right). \diamond$$

After typing up this TRIBONACCI task, now I find the following webpage:

[https://en.wikipedia.org/wiki/Generalizations\\_of\\_Fibonacci\\_numbers#Tribonacci\\_numbers](https://en.wikipedia.org/wiki/Generalizations_of_Fibonacci_numbers#Tribonacci_numbers)

## Number theory of Fib sequence

T.fol problem and soln is from [StackExchange](#).

**13: Five-Fib thm.** For prime  $p$ , fibonacci number  $f_{p-1}$  is divisible by  $p$  IFF  $p \equiv_5 \pm 1$ .  $\diamond$

*Note.* The result holds for  $p=2$ , as both parts fail. Henceforth,  $p$  is an odd prime.  $\square$

**Pf ( $\Leftarrow$ ).** Suppose  $p \equiv_5 +1$  or  $p \equiv_5 -1$ . Then  $p$  is a 5-QR. Courtesy Quadratic Reciprocity, 5 is a  $p$ -QR; this, since 5 is 4POS.

Henceforth,  $\equiv$  means  $\equiv_p$ , and  $\langle \cdot \rangle$  means  $\langle \cdot \rangle_p$ .

Let  $\sigma$  be such that  $\sigma^2 \equiv 5$ ; use  $\hat{\sigma} := \langle 1 \div \sigma \rangle$  for its reciprocal. Let  $\hat{2} := \langle 1 \div 2 \rangle$ . Binet's formula for  $f_n$  is

$$f_n = \frac{1}{\varphi} \cdot [\alpha^n - \beta^n], \quad \text{where} \\ \alpha, \beta = (1 \pm \varphi) \cdot \frac{1}{2}, \quad \text{where } \varphi := \sqrt{5}.$$

**The Idea.** Follow Binet by mimicking  $f_n$  in  $\mathbb{Z}_p$ :

$$\begin{aligned} g_n &:= \hat{\sigma} \cdot \left[ ([1+\sigma] \cdot \hat{2})^n - ([1-\sigma] \cdot \hat{2})^n \right] \\ \text{**;} \quad &= \hat{\sigma} \cdot \left[ (1+\sigma)^n - (1-\sigma)^n \right] \cdot \hat{2}^n. \end{aligned}$$

**Fermat's Little Thm.** There is no oddprime  $p$  for which  $5 \equiv 1$ . Hence the above  $\sigma \not\equiv \pm 1$ . Thus  $1+\sigma$  and  $1-\sigma$  are  $p$ -units, as is  $\hat{2}$ .

We may thus apply FLiT to conclude that

$$\begin{aligned} g_{p-1} &\equiv \hat{\sigma} \cdot \left[ ([1+\sigma] \cdot \hat{2})^{p-1} - ([1-\sigma] \cdot \hat{2})^{p-1} \right] \\ \text{**;} \quad &\stackrel{\text{FLiT}}{=} \hat{\sigma} \cdot [1 - 1] = 0. \end{aligned}$$

**Showing  $g_N \equiv f_N$ .** IStEstablish  $2^N g_N \stackrel{?}{\equiv} 2^N f_N$ , since  $2 \perp p$ . To this end, let  $H$  be the largest integer with  $2H + 1 \leq N$ . Then  $\varphi \cdot 2^N f_N$  equals

$$\begin{aligned} (1+\varphi)^N - (1-\varphi)^N &= \sum_{j=0}^N \binom{N}{j} \cdot [1 - (-1)^j] \varphi^j \\ &= \sum_{d=0}^H \binom{N}{2d+1} \cdot 2 \varphi^{2d+1} \\ &= 2\varphi \cdot \sum_{d=0}^H \binom{N}{2d+1} \cdot 5^d. \end{aligned}$$

Hence  $2^N f_N$  equals

$$\frac{1}{\varphi} \cdot \left[ (1+\varphi)^N - (1-\varphi)^N \right] = 2 \sum_{d=0}^H \binom{N}{2d+1} \cdot 5^d \stackrel{\text{note}}{\equiv} 2N.$$

The same algebra shows  $2^N g_N$  is mod-5 congruent to

$$\hat{\sigma} \cdot \left[ (1+\sigma)^N - (1-\sigma)^N \right] \equiv 2 \sum_{d=0}^H \binom{N}{2d+1} \cdot 5^d \equiv 2N. \quad \blacklozenge$$

**Pf ( $\Rightarrow$ ).** ??  $\blacklozenge$

**13a: Corollary.** For  $n$  a natnum:  $2^n \cdot f_n \equiv_5 2n$ .  $\diamond$