

This IOP is due **2PM, Thur., 20Apr2017**, slid *completely* under my office door, 402 LITTLE HALL. This sheet is “Page 1/N”, and you’ve labeled the rest as “Page 2/N”... “Page N/N”.

**Y1:** Show no work.

**a** Both  $\sim$  and  $\bowtie$  are equiv-relations on a set  $\Omega$ . Define binrels **I** and **U** on  $\Omega$  as follows.

Define  $\omega \mathbf{U} \lambda$  IFF Either  $\omega \sim \lambda$  or  $\omega \bowtie \lambda$  [or both].

Define  $\omega \mathbf{I} \lambda$  IFF Both  $\omega \sim \lambda$  and  $\omega \bowtie \lambda$ .

So “**U** is an equiv-relation” is:  $T \quad F$

So “**I** is an equiv-relation” is:  $T \quad F$

**b** Let  $\mathcal{P}_\infty$  denote the family of all **co-finite** subsets of  $\mathbb{N}$ . That is, a subset  $S \subset \mathbb{N}$  is an *element* of  $\mathcal{P}_\infty$  IFF  $\mathbb{N} \setminus S$  is finite. Define relation  $\bowtie$  on  $\mathcal{P}_\infty$  by:  $A \bowtie B$  IFF  $A \cap B$  is infinite.

Stmt “This  $\bowtie$  is an equivalence-relation” is:  $T \quad F$

**Y3:** [For free: **Union Thm:** A countable union of countable-sets is countable. Also, **Finite-subset Thm:** The collection of **finite** subsets of a countable set, is countable. If needed, use  $\mathcal{P}_{\text{Fin}}(S)$  for the collection of *finite* subsets of a set  $S$ , and use  $\mathcal{P}_\infty(S)$  for the collection of **infinite** subsets of  $S$ .] Below, a **blip** is an *infinite* set of natnums. A **family**,  $\mathcal{F}$ , is a set [not a multiset] of blips, i.e,  $\mathcal{F} \subset \mathcal{P}_\infty(\mathbb{N})$ .

**a** Suppose,  $\forall B, C \in \mathcal{F}$ , that  $[B \neq C] \implies [B \cap C = \emptyset]$ . Construct, with proof, an *injection*  $g: \mathcal{F} \hookrightarrow \mathbb{N}$ , to conclude that every such family,  $\mathcal{F}$ , must only be countable.

**b** Weaken the hypothesis on  $\mathcal{F}$  to: For all  $B, C \in \mathcal{F}$ :  $[B \neq C] \implies |B \cap C| \leq 1$ . Prove  $\mathcal{F}$  is tiny; only countable.

Weaken further to  $[B \neq C] \implies |B \cap C| \leq 2$ , yet still prove  $\mathcal{F}$  countable. Now weaken to  $[B \neq C] \implies |B \cap C| \leq 3$ , and prove  $\mathcal{F}$  is only countable. *Generalize!*

**c** [Challenging/Creative; A converse.] Construct a *specific uncountable* family  $\mathcal{U}$ , so that: For all distinct  $B, C \in \mathcal{U}$ : Intersection  $B \cap C$  is finite.

End of IndividualOP-Y

Your 2 essay(s) must be TYPED, and Double spaced. Use the **Print/Revise** cycle to produce good, well thought out, essays. Start each essay on a new sheet. Do not restate the problem; just solve it.

**Y2:** [Here, “graph” means “non-void finite simple graph”.]

A graph  $M$  is **gluing-good** if, for all graphs  $H, S$  having  $M$  as a subgraph, necessarily

$$\dagger: \quad \mathcal{P}_G(x) = \frac{\mathcal{P}_H(x) \cdot \mathcal{P}_S(x)}{\mathcal{P}_M(x)},$$

whenever  $G$  is an  $M$ -gluing of  $H$  with  $S$ .

**i** For each posint  $N$ , prove that the complete graph  $K_N$  is gluing-good.

**ii** For  $N = 3, 4, \dots$ , show that the path-graph  $P_N$  is *not* gluing-good, as follows. Exhibit graphs  $H_N$  and  $S_N$ , that can be glued [give a specific gluing] over  $P_N$ , to produce a graph  $G_N$  such that  $\mathcal{P}_{G_N}(x) \neq \frac{\mathcal{P}_{H_N}(x) \cdot \mathcal{P}_{S_N}(x)}{\mathcal{P}_{P_N}(x)}$ .

**iii** Generalize: Prove that if  $M$  is *not* a complete graph, then  $M$  not gluing-good.

**Y1:** \_\_\_\_\_ 30pts

**Y2:** \_\_\_\_\_ 000pts

**Y3:** \_\_\_\_\_ 000pts

**Total:** \_\_\_\_\_ 30pts

**HONOR CODE:** “I have neither requested nor received help on this exam other than from my professor (or his colleague).”  
Name/Signature/Ord

Ord: \_\_\_\_\_

*Folks, I've had a wonderful time Problem-Solving with you. Do consider my Combinatorics course in Fall 2017. Stop by in future semesters for Math/chess/coffee.*

*Cheers, Coun-SELO-r King*