

## Ramsey-like theorems

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ABSTRACT: Ramsey-like theorems. Hales-Jewett.

In this note, a **coloring** is an edge-coloring unless specified otherwise.

**Entrance.** (jk: Jeudi, 26 mars, 1985) Let  $K_n$  denote the complete graph on  $n$  vertices. Let  $\mathcal{R}(n)$  denote the  $n^{\text{th}}$  **Ramsey number**, that is, the smallest  $M$  such that if one colors each edge of  $K_M$  either **blue** or **green**, then there must exist a monochromatic subgraph isomorphic to  $K_n$ .

**Notation.** For vertex  $\mathbf{w}$  in a blue-green graph, let  $bV(\mathbf{w})$  be the set of vertices connected to  $\mathbf{w}$  by a **blue** edge; use  $gV(\mathbf{w})$  for those connected to  $\mathbf{w}$  by **green** edges. We thus have vertex degrees

$$b\text{Deg}(\mathbf{w}) := |bV(\mathbf{w})| \quad \text{and} \quad g\text{Deg}(\mathbf{w}) := |gV(\mathbf{w})|.$$

Hence  $b\text{Deg}(\mathbf{w}) + g\text{Deg}(\mathbf{w}) = \text{Deg}(\mathbf{w})$ .

**1a: Theorem.** For  $n = 1, 2, 3, \dots$ ,

$$1a': \quad \mathcal{R}(n+1) \leq 2[1 + [n-1]\mathcal{R}(n)]. \quad \diamond$$

**Proof.** With  $M := \text{RhS}(1a')$ , let  $\mathbb{V}$  denote the vertex set of  $K_M$ . FT SOC, assume we have edge-colored  $K_M$  so that it has no monochromatic copy of  $K_{n+1}$ .

Pick a  $\mathbf{u}_0 \in \mathbb{V}$ . WLOG generality at least half of the edges from  $\mathbf{u}_0$  are **blue**. Letting  $B_0 := b\text{Deg}(\mathbf{u}_0)$ , we have

$$*: \quad B_0 \geq \left\lceil \frac{M-1}{2} \right\rceil \stackrel{\text{note}}{=} 1 + [n-1] \cdot \mathcal{R}(n).$$

**Building a blue complete-subgraph.** Pick a vertex  $\mathbf{u}_1 \in bV(\mathbf{u}_0)$ . The set  $bV(\mathbf{u}_0) \setminus \{\mathbf{u}_1\}$  can have at most  $\mathcal{R}(n) - 1$  vertices not belonging to  $bV(\mathbf{u}_1)$ . If otherwise, then there would be a copy of  $K_{\mathcal{R}(n)}$  in our graph, disjoint from  $\mathbf{u}_0$  and  $\mathbf{u}_1$ , such that each vertex in this copy was connected to  $\mathbf{u}_0$  by a **blue** edge and to  $\mathbf{u}_1$  by a **green** edge. This copy would, by hypothesis, contain a monochromatic copy –call it  $H$ – of  $K_n$ . Were  $H$  **blue**, then  $H \sqcup \{\mathbf{u}_0\}$  is a **blue**  $K_{n+1}$ . Else,  $H$  is **green**, so  $H \sqcup \{\mathbf{u}_1\}$  is a **green**  $K_{n+1}$ . Either is a  $\otimes$ .

Consequently the set  $bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1)$  has at least

$$|bV(\mathbf{u}_0)| - |\{\mathbf{u}_1\}| - [\mathcal{R}(n) - 1] \stackrel{\text{note}}{=} B_0 - \mathcal{R}(n)$$

many vertices. Pick some  $\mathbf{u}_2$  in  $bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1)$ . The same argument shows  $bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1) \cap bV(\mathbf{u}_2)$  has at least

$$[B_0 - \mathcal{R}(n)] - \mathcal{R}(n)$$

many vertices.

Continue. At stage  $n-1$  we will have chosen distinct vertices  $\mathbf{u}_0, \mathbf{u}_1, \dots, \mathbf{u}_{n-1}$ . Further, the intersection of their **blue** vertex-sets will satisfy

$$|bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1) \cap \dots \cap bV(\mathbf{u}_{n-1})| \geq B_0 - [n-1]\mathcal{R}(n)$$

and hence be non-empty, by (\*). Picking a vertex  $\mathbf{u}_n$  in this intersection, now every pair of vertices in  $\{\mathbf{u}_0, \mathbf{u}_1, \dots, \mathbf{u}_n\}$  is connected by a **blue** edge.  $\otimes \blacklozenge$

**Remark.** Values  $\mathcal{R}(1)$ ,  $\mathcal{R}(2)$  and  $\mathcal{R}(3)$  are 1, 2 and 6 respectively. Using the above theorem, easily

$$\mathcal{R}(n+1) \leq [n! \cdot 2^n]. \quad \square$$

**More colors.** If we allow  $\mu \in \mathbb{Z}_+$  many colors, then we get a corresponding sequence of Ramsey numbers  $\mathcal{R}_\mu(n)$  for  $n = 1, 2, \dots$ . Here,  $\mathcal{R}_1(n) = n$ , and  $\mathcal{R}_2(\cdot)$  is another name for  $\mathcal{R}(\cdot)$ .

For each  $\mu > 1$ , lump the  $\mu$  colors into two **super-colors** consisting each of at most  $h := \lceil \mu/2 \rceil$  colors. Applying  $\mathcal{R}(\cdot)$  to these two supercolors, implies that

$$\mathcal{R}_\mu(n) \leq \mathcal{R}(\mathcal{R}_h(n)) \leq \mathcal{R}(\mathcal{R}(\dots \mathcal{R}(n) \dots)),$$

where the RhS has  $\lceil \log_2(\mu) \rceil$  occurrences of  $\mathcal{R}(\cdot)$ .  $\square$

### Infinite Ramsey's Theorem

(jk, Feb1989) Let  $K_\infty$  denote the complete graph on denumerably many vertices.

**1b: Infinite Ramsey's Theorem.** Each edge  $\overline{uv}$  of  $H = (\mathbb{V}, \mathbb{E})$ , a  $K_\infty$  graph, is colored either blue or green. Then  $H$  includes a monochromatic  $K_\infty$ .  $\diamond$

*Proof.* Let  $W_0 := \mathbb{V}$ . At stage  $N$  we have vertex-sets

$$W_0 \supset W_1 \supset \dots \supset W_j \supset \dots \supset W_N,$$

each infinite. Moreover, we have vertices  $(\mathbf{u}_j)_{j=0}^{N-1}$ , with  $\mathbf{u}_j \in W_j \setminus W_{j+1}$ , so that each edge-set

$$\mathbb{E}_j := \{\overline{v\mathbf{u}_j} \mid v \in W_{j+1}\}$$

is monochromatic.

Continue the induction by picking an arbitrary vertex  $\mathbf{u}_N \in W_N$ . Since  $W_N \setminus \{\mathbf{u}_N\}$  is infinite, it includes an infinite set of vertices  $\mathbf{v}$  so that  $\{\overline{v\mathbf{u}_N}\}_v$  is monochromatic. Define  $W_{N+1}$  to be this infinite set of vertices.

**An infinite monochromatic sequence.** Each of the sets  $\mathbb{E}_0, \mathbb{E}_1, \dots$  is either blue or green. Thus there is a subsequence  $N_1 < N_2 < \dots$  so that, WLOG, each  $\mathbb{E}_{N_j}$  is green. The consequence is that the complete graph with vertex-set  $(\mathbf{u}_{N_j})_{j=1}^\infty$  has every edge green.  $\blacklozenge$

**1c: Observation.** Infinite Ramsey's Thm implies the Finite Ramsey's Thm, by a compactness argument.  $\diamond$

*Proof.* Consider  $K_\infty$ , the complete graph with vertex-set  $\mathbb{V} = \{1, 2, \dots\}$ .

FTSOC suppose the, say, 5<sup>th</sup> Ramsey number were not finite. Then for each  $n$ , there would exist a blue-green coloring of the edges on vertex-set  $\{1, \dots, n\}$  –call this colored graph  $H_n$ – so that  $H_n$  does not include a monochromatic copy of  $K_5$ .

Interpret  $H_n$  as an edge-coloring of  $K_\infty$ : Color plum each edge that is not between vertices  $[1..n]$ . (I.e, for posints  $\mathbf{u} < \mathbf{v}$ : If  $\mathbf{v} > n$  then color edge  $\overline{u\mathbf{v}}$  plum.)

We now have a sequence  $H_1, H_2, H_3, \dots$  of  $\{\text{blue, green, plum}\}$  colorings of  $K_\infty$ . Since each edge

of  $K_\infty$  has only finitely many [three] possible colors, the sequence of colorings has a convergent subsequence  $(H_{n_j})_{j=1}^\infty$ . The coloring obtained in the limit, an edge-coloring of  $K_\infty$ , has *no plum*. But  $\infty$ -Ramsey's-Thm guarantees a monochromatic copy of  $K_\infty$ ; WLOG it is green. Letting  $\mathbf{u}_1 < \mathbf{u}_2 < \dots < \mathbf{u}_5$  be its first 5 vertices, this green  $K_5$  was *already* a subgraph of  $H_{\mathbf{u}_5}$ .  $\otimes$

**2a: Two-variable  $\mathcal{R}()$ .** For posints  $\mathbf{b}, \mathbf{g}$ , let  $\mathcal{R}(\mathbf{b}, \mathbf{g})$  be the smallest  $M$  st. each blue-green-coloring of  $K_M$  has either a blue  $K_b$  or a green  $K_g$ . So  $\mathcal{R}(n) = \mathcal{R}(n, n)$ .

Evidently,  $\mathcal{R}(\mathbf{b}, 1) = 1$  and  $\mathcal{R}(\mathbf{b}, 2) = \mathbf{b}$ ; and this holds symmetrically, since  $\mathcal{R}(\mathbf{b}, \mathbf{g}) = \mathcal{R}(\mathbf{g}, \mathbf{b})$ .  $\square$

**2b: Theorem.** For all  $\mathbf{b}, \mathbf{g} \in \mathbb{Z}_+$ ,

$$*: \quad \mathcal{R}(\mathbf{b}, \mathbf{g}) \leq \mathcal{R}(\mathbf{b}-1, \mathbf{g}) + \mathcal{R}(\mathbf{b}, \mathbf{g}-1). \quad \diamond$$

*Proof.* Let  $M := \text{RhS}(*)$ . Fix a vertex  $\mathbf{w}$ ; it has  $M-1$  edges, so has either at least  $\mathcal{R}(\mathbf{b}-1, \mathbf{g})$  blue-edges, or at least  $\mathcal{R}(\mathbf{b}, \mathbf{g}-1)$  green-edges. WLOG

$$|bV(\mathbf{w})| \geq \mathcal{R}(\mathbf{b}-1, \mathbf{g}). \quad \left[ \text{Recall } bV(\mathbf{w}) \text{ is the set of vertices blue-connected to } \mathbf{w}. \right]$$

If the graph induced by  $bV(\mathbf{w})$  has a green  $K_g$ , then DONE. Else, the induced graph admits a blue  $K_{\mathbf{b}-1}$  which, together with vertex  $\mathbf{w}$ , induces a blue  $K_b$ .  $\blacklozenge$

**2c: Exercise:** Coloring with  $\mu$  colors, let  $\mathcal{R}(n_1, \dots, n_\mu)$  be the smallest  $M$  st. each  $\mu$ -coloring,  $H$ , of  $K_M$  has an index  $j$  for which  $H$  admits a  $K_{n_j}$  with all edges the  $j^{\text{th}}$ -color. *Exer-E1: What is the  $\mu$ -color analog of Thm 2b?* [Hint: Don't jump to conclusions.]  $\square$

**Examples from Bona's text**

In class we proved  $\mathcal{R}(3)=6$ . In this instance (2b\*) is sharp, as

$$\mathcal{R}(3,3) \leq 2 \cdot \mathcal{R}(3,2) = 2 \cdot 3 = 6.$$

Using Thm 2b again,

$$\mathcal{R}(3,4) \leq \mathcal{R}(2,4) + \mathcal{R}(3,3) = 4 + 6 = 10.$$

**3a: Claim:**  $\mathcal{R}(3,4) = 9$ . ♦

*Pf of  $\mathcal{R}(3,4) \leq 9$ .* FTSOC, suppose  $H$  is a blue-green-coloring of  $K_9$  with no blue triangle, nor green  $K_4$ .

Could *some* vertex  $\mathbf{w}$  have  $\text{bDeg}(\mathbf{w}) \geq 4$ ? Since no blue triangle, no pair of vertices in  $bV(\mathbf{w})$  has a blue-edge; but then  $bV(\mathbf{w})$  induces a green  $K_4$ . ✘

Could *every*  $\mathbf{w}$  have  $\text{bDeg}(\mathbf{w}) = 3$ ? Then the blue-degree-sum is  $3 \cdot 9$ , which is odd. But this degree-sum must also equal twice the number of blue edges. ✘

So *there exists* a vertex  $\mathbf{w}$  with  $\text{bDeg}(\mathbf{w}) \leq 2$ , hence

$$|gV(\mathbf{w})| \geq 8 - 2 = 6.$$

Since  $\mathcal{R}(3) = 6$ , and  $gV(\mathbf{w})$  cannot have a blue triangle, it must have a green triangle,  $G$ . But then  $G \sqcup \{\mathbf{w}\}$  is a green  $K_4$ . ♦

*Pf of  $\mathcal{R}(3,4) > 8$ .* On vertex-set  $\mathbb{V} := [0..8)$ , let  $\equiv$  and  $\oplus$  and  $\ominus$  each operate mod 8.

Connect vertices  $j, k \in \mathbb{V}$  by blue IFF  $k \ominus j$  is either

- 3, producing a blue octagon, **or**
- 4, producing four center-crossing blue edges. ♦

But no triple of numbers from  $\{3, 4\}$  has  $\oplus$ -sum equal to zero. Hence there is *no* blue triangle.

Color green the remaining edges, i.e  $k \ominus j \in \{1, 2\}$ . Difference=2 makes green square  $\{0, 2, 4, 6\}$ , and square  $\{1, 3, 5, 7\}$ . And difference=1 make a green octagon  $\{0, 1, 2, 3, 4, 5, 6, 7\}$ .

FTSOCcontradiction, suppose  $G$  is a green  $K_4$  subgraph. It can't have three vertices in one square [diff=4], so  $G$  has *two* vertices in *each* square, say  $\{0, 2\}$  and  $\{\mathbf{v}, \mathbf{w}\}$ . But  $\mathbf{v}$  must differ from each of 0,2 by 1 or 2; so  $\mathbf{v} \stackrel{\text{must}}{=} 1$ . Ditto,  $\mathbf{w} = 1$ . ✘

The result now allows, courtesy Thm 2b, that

$$\mathcal{R}(4,4) \leq 2 \cdot \mathcal{R}(3,4) = 18.$$

**3b: Claim:**  $\mathcal{R}(4,4) = 18$ . ♦

*Proof.* To show  $\mathcal{R}(4,4) > 17$ , we exhibit a blue-green-coloring of  $K_{17}$  with no monochromatic  $K_4$ .

On vertex-set  $\mathbb{V} := \{0, 1, \dots, 8, -8, -7, \dots, -1\}$ , let  $\equiv$ ,  $\oplus$ ,  $\ominus$  and  $\langle \cdot \rangle$  each operate mod 17. Blue-connect  $j, k \in \mathbb{V}$  IFF  $k \ominus j$  is a mod-17 QR [Quadratic Residue].

	$x$	$\langle x^2 \rangle$	$x$	$\langle x^2 \rangle$
3c:	$\pm 1$	1	$\pm 5$	8
	$\pm 2$	4	$\pm 6$	2
	$\pm 3$	-8	$\pm 7$	-2
	$\pm 4$	-1	$\pm 8$	-4

Thus QR =  $\{\pm 1, \pm 2, \pm 4, \pm 8\}$ ; so our edge-lengths are  $\{1, 2, 4, 8\}$ . And NQR =  $\{\pm 3, \pm 5, \pm 6, \pm 7\}$ , giving rise to the green edges. Multiplying the vertices by a non-QR element, will exchange the blue and green edges [using that mult distributes-over addition]. So: *It suffices to show that there is no blue  $K_4$ .*

**Number Thy.** FTSOC, suppose  $\mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{d} \in \mathbb{V}$  are the vertices of a blue  $K_4$ . WLOG  $\mathbf{a}=0$ ; replace vertex  $x$  by  $x \ominus \mathbf{a}$ . WLOG  $\mathbf{b}=1$ ; replace each  $x$  by  $\langle x/\mathbf{b} \rangle$ . [The mapping  $x \mapsto \langle x/\mathbf{b} \rangle$  is color-preserving, since  $\mathbf{b}$  is a QR.]

The vertices are now  $\{0, 1, \mathbf{c}, \mathbf{d}\}$ .

Vertex  $\mathbf{c}$  must differ from 0 and 1 by values in  $\{1, 2, 4, 8\}$ ; so  $\mathbf{c}$  is 2 or -1; WLOG  $\mathbf{c}=2$ . Vertex  $\mathbf{d}$ , relative to vertices 1 and 2, must be 3 or 0; but 0 is taken. Our supposedly-blue  $K_4$  is  $\{0, 1, 2, 3\}$ .

Alas, edge  $0\bar{3}$  is green. ✘ ♦

## van der Waerden's thm

A *shift*  $\mathbf{s} \in \mathbb{Z}$  and a *gap*  $\mathbf{g} \in \mathbb{Z}_+$ , give rise to *finite* and *infinite arithmetic progressions*:

$$\begin{aligned} \mathbb{A}_n(\mathbf{s}, \mathbf{g}) &:= \{\mathbf{s} + j\mathbf{g} \mid j \in [0..n)\}; \\ \mathbb{A}(\mathbf{s}, \mathbf{g}) = \mathbb{A}_\infty(\mathbf{s}, \mathbf{g}) &:= \{\mathbf{s} + j\mathbf{g} \mid j \in \mathbb{Z}\}. \end{aligned}$$

Use  $\mathbb{A}_n$  to refer to *some*  $n$ -term A.P, when the shift and gap are irrelevant.

**4a: van der Waerden theorem.** *For each posint  $n$ , there is a finite  $M$  so that each 2-coloring (of the elements) of  $\mathbb{A}_M$ , produces a monochromatic  $\mathbb{A}_n$ .*  $\diamond$

The smallest such  $M$  is written  $\mathcal{W}(n)$ . For example,  $\mathcal{W}(3) = 9$ . If  $\mu$  many colors are allowed, then we write  $\mathcal{W}_\mu(n)$  for the minimum  $M$ .

The infinite analog of vdW thm fails, as shown next.

**4b: Obs.** *There exists a 2-coloring of  $\mathbb{Z}$  which admits no monochromatic sub- $\mathbb{A}_\infty$ .*  $\diamond$

**Proof.** We produce a coloring for which the only monochromatic arithmetic-progs are finite.

Take a surjection  $n \mapsto (\mathbf{s}_n, \mathbf{g}_n)$  from  $\mathbb{Z}_+ \rightarrow \mathbb{Z} \times \mathbb{Z}_+$ , and let  $\mathcal{A}(n) := \mathbb{A}_\infty(\mathbf{s}_n, \mathbf{g}_n)$ .

We color  $\mathbb{Z}$  in stages. At stage  $n$ : Paint one **blue** and one **green**, the two posints in  $\mathcal{A}(n)$  which are closest to zero, and were not yet colored. Then paint one **blue** and one **green**, the two negints in  $\mathcal{A}(n)$  which are closest to zero, and were not yet colored.

Send  $n \nearrow \infty$ . Finally, color arbitrarily the remaining uncolored integers.  $\blacklozenge$

### Hypergraphs

Let  $\mathcal{P}_4(\mathbb{V})$  be the collection of cardinality-4 subsets of  $\mathbb{V}$ ; so it has  $\binom{|\mathbb{V}|}{4}$  many members. A “4-**hypergraph** on  $\mathbb{V}$ ” is a pair  $H := (\mathbb{V}, \mathbb{E})$  where  $\mathbb{E} \subset \mathcal{P}_4(\mathbb{V})$  is the set of **hyper-edges** of  $H$ . This cardinality, 4, is called the **rank** of  $H$ .

A subset  $V' \subset \mathbb{V}$  determines a sub-hypergraph  $H' := (V', \mathbb{E}')$  of  $H$ , where

$$\mathbb{E}': \quad \mathbb{E}' := \{S \in \mathbb{E} \mid S \subset V'\}.$$

By the way, the “*complete* 4-hypergraph on  $\mathbb{V}$ ” is  $(\mathbb{V}, \mathcal{P}_4(\mathbb{V}))$ . Setting  $M := |\mathbb{V}|$ , we’ll use  $K_M^{(4)}$  to refer to this hypergraph. [So  $K_M^{(2)}$  is our usual  $K_M$ .]

5a: *Defn.* Fix a colorset  $\mathcal{C} := \{\text{blue}, \text{green}\}$ .

#### A $\mathcal{C}$ -coloring of $H$

is a map  $\mathbf{f}: \mathbb{E} \rightarrow \mathcal{C}$ . Given a subset  $V' \subset \mathbb{V}$ ,

let  $\mathbf{f}|_{V'}$  mean the coloring  $\mathbf{f}|_{\mathbb{E}'}$  defined by (5).

Consider a rank  $\rho \in [2.. \infty)$  and posint  $n$ . Suppose there is a posint  $M$  st:

For each  $\mathcal{C}$ -coloring of the complete hypergraph  $K_M^{(\rho)}$ , there is a cardinality- $n$  subset  $W \subset [1.. M]$  so that coloring  $\mathbf{f}|_W$  is constant.

(IOWords, this edge-colored  $K_M^{(\rho)}$  admits a monochromatic  $K_n^{(\rho)}$ .) The smallest such  $M$  is the **hypergraph Ramsey number**  $R^{(\rho)}(n)$ .  $\square$

5b: **Hypergraph Ramsey Thm.** Fix a rank  $\rho \in \mathbb{Z}_+$ . Then

- i: Each coloring of  $K_\infty^{(\rho)}$  admits a monochromatic  $K_\infty^{(\rho)}$ -subgraph.
- ii: Each posint  $n$ : Ramsey number  $R^{(\rho)}(n)$  is finite.  $\diamond$

*Proof of (ii).* The analogous compactness argument of (1c) works here.  $\diamond$

*Proof of (i).* We induct on  $\rho$ . We’ll show the induction for  $\rho=5$ , assuming the  $\rho=4$  case.

Our vertex set is  $\mathbb{V} := \{1, 2, \dots\}$ , and we are given a coloring  $\mathbf{f}: \mathcal{P}_5(\mathbb{V}) \rightarrow \mathcal{C}$ . Suppose we could find an infinite subset  $\mathcal{W} \subset \mathbb{V}$  and a color map  $g: \mathcal{P}_4(\mathcal{W}) \rightarrow \mathcal{C}$  with this property:

†: For each  $S \in \mathcal{P}_4(\mathcal{W})$  and each  $y \in \mathcal{W}$  with  $y > \text{Max}(S)$ , the color  $\mathbf{f}(S \sqcup \{y\})$  equals  $g(S)$ .

The rank=4 case of (5b) asserts there is an  $\infty$ -subset  $\mathbf{X} \subset \mathcal{W}$  so that

Our  $g$ -coloring is constant on  $\mathcal{P}_4(\mathbf{X})$ ; say *blue*.

Given a  $T \in \mathcal{P}_5(\mathbf{X})$ , write it as  $T = \{w_1, \dots, w_4, w_5\}$  with  $w_1 < \dots < w_5$ . By (†), then,

$$\mathbf{f}(T) = g(\{w_1, \dots, w_4\}) = \text{blue}.$$

Hence  $\mathbf{f}$  is constant *blue* on  $\mathcal{P}_5(\mathbf{X})$ .

**Building  $\mathcal{W}$ .** We’ll inductively construct vertices  $w_1 < w_2 < \dots$  and infinite  $\mathbb{V}$ -subsets  $\mathbf{Y}_1 \supset \mathbf{Y}_2 \supset \dots$ . Our  $\mathcal{W}$  will be  $\{w_1, w_2, w_3, \dots\}$ .

Let  $w_1 := 1$ ,  $\mathbf{I}_1 := \{w_1\}$  and  $\mathbf{Y}_1 := [2.. \infty)$ .

At **STAGE  $k$** : We have  $\mathbf{I}_k := \{w_1, \dots, w_k\}$ , and a partially-defined  $g()$ , defined on  $\mathcal{P}_4(\mathbf{I}_k)$ . We have an infinite vertex-set  $\mathbf{Y}_k$ , such that:

- i: Our  $w_k < y$ , for each  $y \in \mathbf{Y}_k$ .
- ii: For each  $S \in \mathcal{P}_4(\mathbf{I}_k)$ , and each  $y \in \mathbf{Y}_k$ , the color  $\mathbf{f}(S \sqcup \{y\})$  equals  $g(S)$ .

For **STAGE  $[k+1]$** , define  $w_{k+1} := \text{Min}(\mathbf{Y}_k)$ , and temporary set

$$J_0 := \mathbf{Y}_k \setminus \{w_{k+1}\}.$$

Let  $S_1, S_2, \dots, S_L$  be some enumeration of those cardinality-4 subsets of  $\mathbf{I}_{k+1}$  that own  $w_{k+1}$ .

There is a color, say, *blue*, and an infinite set of  $y \in J_0$ , so that  $\mathbf{f}(S_1 \sqcup \{y\})$  is *blue*. Extend  $g()$  by defining  $g(S_1) := \text{blue}$ . Use  $J_1$  for this set of points  $y$ .

There is a color, say, *green*, and an infinite set of  $y \in J_1$ , so that  $\mathbf{f}(S_2 \sqcup \{y\})$  is *green*. Define  $g(S_2) := \text{green}$ . Use  $J_2$  for this set of points  $y$ .

*Continue*, until you have shrunk to  $J_L$ . Lastly, let  $\mathbf{Y}_{k+1} := J_L$ .  $\diamond$

**5c: Defn.** The **Erdős-Szekeres number**  $ES(n)$ , is the smallest posint  $M$  so that *each* collection of  $M$  points in the plane with no three colinear, has a subset of  $n$  points which form a *convex*  $n$ -gon. [Caveat: There are at least two different results called the Erdős-Szekeres thm.]  $\square$

**5d: ES-Theorem.** For the Erdős-Szekeres number,

$$ES(n) \leq R^{(3)}(n, n) =: M. \quad \diamond$$

**Proof.** With vertex-set  $[1..M]$ , construct a **Cyan-Amber**-coloring of  $K_{[1..M]}^{(3)}$ , as follows: For each triple  $\mathbf{u} < \mathbf{v} < \mathbf{w}$  in  $[1..M]$ , color the  $\{\mathbf{u}, \mathbf{v}, \mathbf{w}\}$ -edge **Cyan** if the  $\mathbf{u} \rightarrow \mathbf{v} \rightarrow \mathbf{w} \rightarrow \mathbf{u}$  traversal is ClockWise [**CW**]; otherwise, paint **Amber** the  $\{\mathbf{u}, \mathbf{v}, \mathbf{w}\}$ -edge, since its traversal is Anti-clockWise [**AW**].

By hypothesis, there exists an  $n$ -set  $S \subset [1..M]$ , so that all the  $K_S^{(3)}$ -edges are, say, **CW**.

**Convex  $n$ -gon.** FTSOC, suppose the  $n$  points of  $S$  do *not* form a convex  $n$ -gon. Then some point  $P \in S$  is in the convex-hull of  $S$ . So there are distinct points  $\mathbf{u} < \mathbf{v} < \mathbf{w}$  in  $S$ , with an  $S$ -point  $P \in \text{Hull}(\{\mathbf{u}, \mathbf{v}, \mathbf{w}\})$ .

Recall  $\mathbf{u} \rightarrow \mathbf{v} \rightarrow \mathbf{w}$  is **CW**. We must have  $P > \mathbf{u}$ ; else  $P \rightarrow \mathbf{u} \rightarrow \mathbf{w}$  is **AW**.

And  $P \overset{\text{must}}{>} \mathbf{v}$ ; else  $\mathbf{u} \rightarrow P \rightarrow \mathbf{v}$  is **AW**. Continuing,  $P \overset{\text{must}}{>} \mathbf{w}$ ; else  $\mathbf{v} \rightarrow P \rightarrow \mathbf{w}$  is **AW**. But now,  $\mathbf{u} \rightarrow \mathbf{w} \rightarrow P$  is **AW**.  $\otimes$   $\blacklozenge$

## §A Hales-Jewett

### The Statement

Suppose we have a finite alphabet  $\mathbb{A} = \{\mathbf{a}, \mathbf{b}, \dots\}$  and we fix a length  $\mathbf{h}$ . A *degree- $\mathfrak{D}$  polynomial*  $f(x_1, \dots, x_{\mathfrak{D}})$  over  $\mathbb{A}^{\times \mathbf{h}}$  is a word

$$6: \quad f \in [\mathbb{A} \sqcup \{x_1, \dots, x_{\mathfrak{D}}\}]^{\times \mathbf{h}}$$

where each variable  $x_j$  occurs at least once in  $f$ . We evaluate  $f()$  at a  $\mathfrak{D}$ -tuple of  $\mathbb{A}$ -letters by plugging them in for the  $\mathfrak{D}$  variables. The range of this polynomial is a subset of  $\mathbb{A}^{\times \mathbf{h}}$  and has  $|\mathbb{A}|^{\mathfrak{D}}$  members. This range is called a  *$\mathfrak{D}$ -dimensional (affine) subspace*. The RHS of (6) implies

$$7: \quad \text{There are at most } [|\mathbb{A}| + \mathfrak{D}]^{\mathbf{h}} \text{ many } \mathfrak{D}\text{-dimensional subspaces of } \mathbb{A}^{\times \mathbf{h}}.$$

By the way, a 1-dimensional subspace is also called an (affine) *line*.

The Hales-Jewett theorem states that given an *alphabet size*  $\alpha := |\mathbb{A}|$ , a number  $\mu$  of colors and a dimension  $\mathfrak{D}$ :

*There is a function  $\mathbf{h} = \mathbf{h}(\mathfrak{D}, \mu, \alpha)$  so that each  $\mu$ -coloring (coloring by  $\mu$  many colors) of the set of words  $\mathbb{A}^{\times \mathbf{h}}$  will have a monochromatic  $\mathfrak{D}$ -dimensional subspace.*

One cannot guarantee the stronger statement that there is a monochromatic subspace parallel to the coordinate axes ie. where each variable in the word of (6) occurs exactly once. This is already false in the  $\mathfrak{D}=1$  case: Let the *color* of a word in  $\{0, 1\}^{\times \mathbf{h}}$  be the mod-2 sum of its bits. Then a *line* consists of a pair of  $\mathbf{h}$ -words  $u0w$  and  $u1w$  differing in a single bit-position—which therefore have different colors.